

On colorful edge triples in edge-colored complete graphs

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Abstract

An edge-coloring of the complete graph K_n we call F -caring if it leaves no F -subgraph of K_n monochromatic and at the same time every subset of $|V(F)|$ vertices contains in it at least one completely multicolored version of F . For the first two meaningful cases, when $F = K_{1,3}$ and $F = P_4$ we determine for infinitely many n the minimum number of colors needed for an F -caring edge-coloring of K_n . An explicit family of $2\lceil \log_2 n \rceil$ 3-edge-colorings of K_n so that every quadruple of its vertices contains a totally multicolored P_4 in at least one of them is also presented. Investigating related Ramsey-type problems we also show that the Shannon (OR-)capacity of the Grötzsch graph is strictly larger than that of the five length cycle.

1 Introduction

Erdős [12] formulated the following Ramsey type problem: What is the least number $f(n, p, q)$ of colors needed for an edge-coloring of the complete graph K_n if it has the property that every subset of p vertices spans at least q distinct colors? Initial investigations by Elekes, Erdős, and Füredi reported in [12] were followed by a more systematic study by Erdős and Gyárfás in [13]. A more general variant of the problem was studied by Axenovich, Füredi, and Mubayi [3], where they color the edges of a not necessarily complete graph G and require that every copy of another graph $H \subseteq G$ receive at least a given number of colors. Axenovich and Iverson [4] investigated the mixed anti-Ramsey numbers that are the maximum and minimum numbers of colors to be used in an edge-coloring of K_n that avoids both monochromatic copies of a fixed graph G and also totally multicolored copies of another fixed graph H .

Here we initiate the investigation of yet another variant that will also be connected to some of the above versions of the classical Ramsey problem. Let $g(n, F)$ denote the minimum number of colors needed in an edge-coloring of K_n if it contains no monochromatic copy of F , furthermore, it contains at least one totally multicolored copy of F on every subset of $|V(F)|$ vertices.

If F has only two edges then the two conditions coincide and we get well-known numbers: for F being two incident edges we need a proper edge coloring of K_n , while in case of two non-incident edges we ask the chromatic number of the complementary graph of the line graph of K_n , or equivalently, the chromatic number of the special Kneser graph $KG(n, 2)$ (cf. [24]), which is $n - 2$. The first meaningful cases of our problem thus belong to graphs F with 3 edges and 4 vertices, since if F is a triangle, or a complete graph, in general, then we are back to the trivial problem of determining $f(n, p, \binom{p}{2}) = \binom{n}{2}$, with $p \geq 3$, cf. [13]. Thus the first interesting cases belong to $F = K_{1,3}$, and $F = P_4$, the star and the path on four vertices. We will determine $g(n, K_{1,3})$ and $g(n, P_4)$ for infinitely many n in Section 2.

Completely multicolored edge triples are investigated in a somewhat different context in [23] and also in the subsequent paper [5]. There we have only three colors but we can color in several rounds. Vera T. Sós asked in 1991, how many 3-edge-colorings of K_n ($n \geq 3$) are needed to have every triangle 3-colored in at least one of them. This question was related to a special case of the so-called perfect hashing problem, cf. [16, 20, 21]. In [23] we gave a construction proving that the requested minimum number is at most $\lceil \log_2 n \rceil - 1$, while an easy argument shows that it is bounded from below by $\lceil \log_3(n - 1) \rceil$. Different edge triples were also considered in [23], but we did not find meaningful bounds in the case when these edges form a P_4 . In Section 3 we give upper bounds for the number of 3-edge-colorings needed if we want that at least one P_4 gets 3-colored on every 4 vertices in at least one of our colorings.

Note that if we wanted to only 2-color every triangle in several rounds of 3-edge-colorings then we would need exactly $\lceil \log_3 k(n) \rceil$ rounds, where $k(n) = \min\{k : R(3; k) > n\}$, where $R(3; k)$ is the minimum size of a complete graph that cannot be k -edge colored

without creating a monochromatic triangle, i.e., the Ramsey number $R(3, 3, \dots, 3)$ with the number of 3's being k . (We simply have to encode each color with a ternary sequence to convert the coloring into $\lceil \log_3 k(n) \rceil$ 3-edge-colorings.) The growth rate of this number is well-known to be closely related to the Shannon (OR-)capacity of triangle-free graphs. Elaborating on this connection we observe, using a result of Cropper, Gyárfás, and Lehel [9], that the answer to a famous question about the growth rate of $R(3 : k)$ depends on whether the Shannon (OR-)capacity of Mycielski graphs go to infinity. As a first modest step towards such investigations we show that the Shannon (OR-)capacity of the Grötzsch graph is strictly more than that of C_5 , the cycle of length five. These questions are discussed in the last subsection of Section 3.

2 $K_{1,3}$ - and P_4 -caring colorings

Along with introducing some other related parameters, let us repeat the definition of $g(n, F)$ that we already defined less formally in the Introduction.

Definition 1. *Let F be any fixed graph and $n \geq |V(F)|$. An edge-coloring of K_n in which no copy of F is monochromatic and at the same time every subset of $|V(F)|$ vertices contains a totally multicolored (rainbow) copy of F is called F -caring. Let $g(n, F)$ denote the minimum number of colors needed to an F -caring coloring of K_n .*

Let $b(n, F)$ denote the minimum number of colors needed for an edge-coloring of K_n that contains no monochromatic copy of F .

Finally, let $a(n, F)$ denote the minimum number of colors needed for an edge-coloring of K_n that makes at least one copy of F totally multicolored in every subset of $|V(F)|$ vertices.

Thus $b(n, F)$ is simply the smallest k for which the Ramsey number $R_k(F, F, \dots, F) > n$, where the index k refers to the number of colors used. It is immediate from the above definition that $g(n, F) \geq \max\{a(n, F), b(n, F)\}$. It is a trivial matter to determine $b(n, K_{1,3})$ exactly, and it is also easy to tell the value of $b(n, P_4)$ for infinitely many values of n .

Proposition 1. *One has $b(n, K_{1,3}) = \lceil \frac{n-1}{2} \rceil$ for every $n \geq 4$.*

Proof. Any color class that contains no $K_{1,3}$ must give a graph with maximum degree at most 2, thus it is a collection of vertex-disjoint paths and cycles. It follows that the number of edges in each color class is at most n and thus the lower bound $b(n, K_{1,3}) \geq \lceil \frac{n-1}{2} \rceil$ follows.

On the other hand, it is well-known that if n is odd, then the edge set of K_n can be partitioned into $\frac{n-1}{2}$ Hamiltonian cycles, see e.g. [7]. Letting the edges of these cycles form the color classes the Proposition follows for odd values of n .

If n is even, then we can simply take an optimal edge-coloring of K_n with $n-1$ colors, and pair up the colors in an arbitrary way. For each pair of colors define a new color

class, obviously, no vertex will connect to three edges in it. This proves the required upper bound for even values of n . \square

For the next proposition we need some notions from the theory of block designs. A Steiner triple system (STS for short) is a family of 3-element subsets of a basic set of n elements having the property that every pair of elements belongs to exactly one triple of the family. An STS is called resolvable if it can be partitioned into subfamilies, called *parallel classes*, in which each element of the basic set belongs to exactly one triple. (That is, in each parallel class the triples are pairwise disjoint and their union is the entire basic set.) Resolvable Steiner triple systems are called Kirkman triple systems (KTS for short) referring to Kirkman's puzzle in the middle of the 19th century asking for the existence of such a system of triples and their parallel classes for $n = 15$, cf. [32]. In full generality it was proven only more than a hundred years later by Ray-Chaudhury and Wilson [29] that Kirkman triple systems exist for all $n \equiv 3 \pmod{6}$. (It is easy to see that they cannot exist for any other n .)

Proposition 2. *If $n \equiv 3 \pmod{6}$, then $b(n, P_4) = \frac{n-1}{2}$.*

Proof. Consider an edge-coloring of K_n with $b(n, P_4)$ colors in which no P_4 is monochromatic. Then each color class is a P_4 -free graph, so it must be a collection of vertex-disjoint triangles and stars. This implies that the number of edges in each color class is at most n . So the number of color classes $b(n, P_4) \geq \frac{n-1}{2}$.

For equality we need a coloring in which all color classes contain only triangles (none of them are stars). As every edge gets exactly one color, this requires that the union of the triangles forming the color classes be a Steiner triple system if we consider the triangles as triples. Furthermore, as each color class should consist of vertex-disjoint triangles, they should form the parallel classes of a Kirkman triple system. Conversely, if a Kirkman triple system on n points exists, then we can use its parallel classes to define color classes of the edges so that each edge is colored with some identifier of the parallel class it belongs to (that is, which contains a triple, that covers this edge). Such a coloring uses $\frac{n-1}{2}$ colors and it lets no P_4 remaining monochromatic. This proves $b(n, P_4) \leq \frac{n-1}{2}$ whenever $n \equiv 3 \pmod{6}$. \square

The two propositions above give obvious lower bounds for the value of $g(n, K_{1,3})$ and $g(n, P_4)$. Below we show that these lower bounds are tight.

Proposition 3. *If $n \equiv 3 \pmod{6}$, then*

$$g(n, K_{1,3}) = \frac{n-1}{2}$$

Proof. Since $g(n, F) \geq b(n, F)$ by definition, $g(n, K_{1,3}) \geq \frac{n-1}{2}$ is immediate from Proposition 1. For the reverse inequality take a coloring based on a Kirkman triple system as

in the proof of Proposition 2, for the numbers in the statement such a system exists by the result of Ray-Chaudhury-Wilson [29]. We show this is a $K_{1,3}$ -caring coloring.

It is clear that no $K_{1,3}$ is monochromatic, since each vertex has degree 2 only within each graph formed by edges of a color class. Note, that if two incident edges belong to the same color class, then the edge extending this pair of edges to a triangle should also belong to the same color class. Now consider a set of four points $\{a, b, c, d\}$. If there is a monochromatic triangle on them formed by, say, vertices a, b, c , then the fourth vertex must be connected to each of a, b , and c in a different color. (If two of them were the same, then it would require an edge of the monochromatic triangle to belong to another monochromatic triangle, which is impossible.) Thus the $K_{1,3}$ spanned by edges da, db, dc is rainbow. If there is no monochromatic triangle on our 4 points, then it cannot contain even two incident edges with the same color and thus all the four $K_{1,3}$'s it contains are rainbow. \square

In what follows we will use the notation $a \oplus_n b$ for the modulo n sum of integers a and b .

Theorem 1. *If n is a power of 3, then*

$$g(n, P_4) = \frac{n-1}{2}$$

Proof. Since $g(n, F) \geq b(n, F)$, we have $g(n, P_4) \geq \frac{n-1}{2}$ from Proposition 2. Now let $n = 3^t$. For the reverse inequality we show a Kirkman triple system that has the property that every quadruple of vertices contains some P_4 the three edges of which belong to three different parallel classes of the KTS. (In what follows we will often refer to the parallel classes as colors. Colorings defined by the parallel classes of a KTS as color classes will be called KTS colorings. Note that no P_4 can be monochromatic in a KTS coloring according to the argument in the proof of Proposition 2).

It is straightforward to check that if a K_4 is edge-colored with at least 4 colors then it must contain a rainbow P_4 . If a K_n is edge-colored with the parallel classes of a KTS then there is only one way not to have 4 differently colored edges on a fixed set of four vertices, namely when it is colored with three colors and the coloring partitions its edge set into three matchings, each containing exactly two edges. (In other words it is an optimal proper edge coloring of this K_4 .) This simply follows from the fact that if two incident edges have the same color in a KTS coloring, then the third edge with which they form a triangle must also belong to the same color class. So it is enough to show a KTS on $n = 3^t$ vertices for every t with the property that the KTS coloring it defines will not partition the edge set of any K_4 into three perfect matchings. Let us call a KTS with this property that it has at least 4 parallel classes intersecting the edge set of any $K_4 \subseteq K_n$ a good KTS.

The smallest non-trivial KTS that exists is on 9 points and is easily shown to have this property as follows. Assume that for some four points $\{a, b, c, d\}$ the coloring it defines colors $\{a, b\}$ and $\{c, d\}$ the same color. Then the edges $\{a, b\}$ and $\{c, d\}$ belong to disjoint

triples of the same parallel class of the KTS. This needs two distinct vertices belonging to those two triples outside the set $\{a, b, c, d\}$. The same can be said about the other two pairs of edges, so altogether we need six vertices outside $\{a, b, c, d\}$, but we only have 5 as there are only 9 points altogether. So the KTS on 9 points is a good KTS.

Now we prove the theorem by induction on t . Assume we have a good KTS on 3^{t-1} points, and show that one exists also on 3^t points. (The base case was just verified for $t = 1$. In fact, we could have taken $t = 0$ for the base case as well, and get the KTS on 9 points also by the induction step.) Partition our 3^t points into three equal classes A_j and label them as (ℓ, j) with $\ell \in \{0, 1, 2\}$ referring to their partition class and $j \in \{0, 1, \dots, 3^{t-1} - 1\}$ their position according to an arbitrary labelling within their class. Take three copies of the existing good KTS on 3^{t-1} points with parallel classes $C_1, \dots, C_m, m = \frac{3^{t-1}-1}{2}$ on the three partition classes and consider the union of their i th parallel classes as the i th parallel class of the KTS we are constructing on 3^t points. For $j = 0, 1, \dots, 3^{t-1} - 1$ we define $C_{m+j+1} := \{(0, a), (1, a \oplus_{3^{t-1}} j), (2, a \oplus_{3^{t-1}} 2j)\} : 0 \leq a < 3^{t-1}$. Then $\{C_1, \dots, C_{\frac{3^t-1}{2}}\}$ is a Kirkman triple system, we show it is good.

Call a K_4 bad (with respect to a KTS) if its edges belong to only 3 parallel classes of the KTS. Consider any four points, we want to show that they cannot induce a bad K_4 (with respect to the KTS just defined). If the four points are all in the same class A_j then they could induce a bad K_4 only if the KTS on 3^{t-1} points we started with would not be a good one, which is not the case. If our four points are in 3 different classes, then exactly one pair of these four points belongs to the same class and thus to a C_k with $0 \leq k \leq m$ with m still meaning the value $\frac{3^{t-1}-1}{2}$. Then this pair of points forms a differently colored edge then the complementary two points in our 4-element set, so the edge set of our K_4 cannot be partitioned into three perfect matchings and thus this K_4 cannot be bad. If the four points are in two classes so, that one is in one class, say A_j , and three of them are in the other class, say A_h then the three edges of our K_4 that has one endpoint in A_j and the other in A_h all belong to different C_k 's with $m < k$ while the remaining three edges all belong to some C_k 's with $k \leq m$. So again, we have at least 4 colors in our K_4 . Lastly, if two of our points are in A_j and two are in A_h with $j \neq h$, then the four edges running between the two classes must belong to at least three different C_k 's with $k > m$. This is because only pairs of independent edges among them can have the same color and if $a, b \in A_j, c, d \in A_h$, and edges $\{a, c\}, \{b, d\}$ have the same color, then we have $c - a \equiv d - b \pmod{3^{t-1}}$. Both $\{a, d\}$ and $\{b, c\}$ should then belong to a different color class, and these two may not be equal, otherwise we also have $c - b \equiv d - a \pmod{3^{t-1}}$, and the last two congruences would imply $b - a \equiv a - b \pmod{3^{t-1}}$ contradicting that a and b are different. The remaining two edges $\{a, b\}$ and $\{c, d\}$ belong to C_k 's with $k \leq m$, so we again have at least four differently colored edges in our K_4 . This completes the proof of the theorem. \square

Remark 1. We have $g(n, F) \geq \max\{a(n, F), b(n, F)\}$ for every graph and we have seen above that for $K_{1,3}$ and P_4 we have $g(n, K_{1,3}) = b(n, K_{1,3})$ and $g(n, P_4) = b(n, P_4)$. In

fact, $a(n, K_{1,3})$ and $a(n, P_4)$ are indeed strictly smaller than $b(n, K_{1,3})$ and $b(n, P_4)$. This can be seen by using that $a(n, K_{1,3}), a(n, P_4) \leq f(n, 4, 4)$ (for the meaning of $f(n, p, q)$ see the beginning of the Introduction), and Erdős and Gyárfás's result from [13] stating that $f(n, 4, 4) \leq O(n^{2/3})$, improved by Mubayi [26] (via an explicit construction) to $f(n, 4, 4) \leq n^{1/2} e^{c\sqrt{\log n}}$ with an absolute constant $c > 0$. Nevertheless, $a(n, F) > b(n, F)$ can also happen. This is trivially so if F is a complete graph, but also for any F with $|E(F)| > \frac{1}{2} \binom{|V(F)|}{2}$ we immediately have $a(n, F) \geq b(n, F)$. For F being a complete graph minus one edge we also trivially have $g(n, f) = a(n, F) = f(n, |V(F)|, \binom{|V(F)|}{2} - 1)$. In particular, $f(n, 4, 5) \geq \frac{5}{6}(n-1)$ is proven in [13] implying that for $F = K_4 \setminus \{e\}$ (where e is an edge of K_4) we have $b(n, F) \leq f(n, 4, 3) < f(n, 4, 5) = a(n, F)$. \diamond

Remark 2. Although an analogous result to that of Ray-Chaudhury and Wilson [29] about the existence of Kirkman triple systems was proven by Hanani, Ray-Chaudhury, and Wilson [18] for quadruple systems (that is the existence of resolvable so-called balanced incomplete block designs of block size 4 whenever a trivial necessary condition is met), it cannot be directly used for generalizing even Proposition 3 to determine $g(n, K_{1,4})$. This is because a colouring based on such designs will not guarantee the 4-coloring of a $K_{1,4}$ subgraph on every 5 vertices. \diamond

3 Coloring in several rounds

3.1 Relaxations of perfect hashing

The perfect hashing problem, investigated in [16, 20, 21] asks for the minimum length of sequences over an alphabet of size $b \geq k$, such that one can give n such sequences so that for any k of them there is a position where all of these k sequences differ. While the answer is trivially $\log_b n$ if $k = 2$, for $k = 3$ the answer is not known even for the smallest meaningful alphabet size $b = k = 3$. Denoting the required minimum length by $t(n, b, k)$ in general, it is known that $t(n, b, k) = \Theta(\log n)$ for any fixed pair (b, k) and the relevant quantity asked for is $c_{b,k} = \liminf_{n \rightarrow \infty} \frac{t(n, b, k)}{\log n}$. The best known bounds on $c_{b,k}$ in general are given in [21], in case of $c_{3,3}$ their values are

$$\frac{1}{\log \frac{3}{2}} \leq c_{3,3} \leq \frac{4}{\log \frac{9}{5}},$$

where, as always in the sequel when not stated otherwise, the logarithms are meant to be on base 2. Numerically the above bounds mean $1.709 \leq c_{3,3} \leq 4.717$.

We remark that the upper bound is implicit in Elias [11], and that for $c_{4,4}$ some improvements were found recently by Dalai, Guruswami, and Radhakrishnan [10]. From now on we are focusing here only on the $b = k = 3$ case.

While this problem is notoriously difficult, it is interesting to look at relaxations where we do not require all triples of our sequences being pairwise different at some position (this is the relation we called "triference" in [23]) but only a selected collection of all

triples should have this property. Such a selected collection is most naturally defined by some underlying structure. This explains what could be the motivation for Vera T. Sós in 1991 for asking how long our sequences should be if the basic set is identified to the edge set of K_n and only those triples should be trifferent at some position that form a triangle. In other words: How many 3-edge-colorings of K_n are needed to make every triangle 3-colored in at least one of them? Though we still do not know the exact answer to this question, it indeed seems to be more easily tractable than the problem of $c_{3,3}$, because here we could at least give an explicit construction (see [23]) that implies a $\lceil \log_2 n \rceil - 1$ upper bound on the length of sequences needed, while an obvious lower bound is $\lceil \log_3(n-1) \rceil$. These two bounds differ first at $n = 9$ for which Blokhuis, Faudree, Gyárfás, and Ruszinkó [5] proved that the upper bound is tight. For a completely solved variant involving orientations in place of colorings, see [19]. (For the general triference problem the best known upper bounds on $c_{3,3}$ are obtained via the probabilistic method, that is, they are not constructive.) If instead of the edges of triangles we want the 3 edges of every $K_{1,3}$ or every 3 independent edges to become trifferent in at least one coloring, then the required minimum number of colorings will be $c_{3,3} \log n + o(\log n)$, see [23]. (The reason is simple: In case of $K_{1,3}$, we have to make all the $(n-1)$ edges incident to a fixed vertex have the property that any three of them is trifferent at some coloring, while encoding the $(n-1)$ or n colors of an optimal proper edge-coloring this way will do job; and a similar reasoning works for three independent edges.)

We did not find non-trivial bounds in [23] for the triples given by the edge sets of P_4 's. A trivial upper bound is $2c_{3,3} \log n + o(\log n)$, since with that many colorings we can make all the $\binom{n}{2}$ edges having the property that any triple is trifferent at some coloring. We also have $\log_3(n-1)$ as a trivial lower bound as any two edges incident to the same fixed vertex should get a different color in some coloring as they appear together in some P_4 . (Note that in this kind of problems it is typically the constant multiplier of $\log n$ that we are seeking for.)

A consequence of Erdős and Gyárfás's [13] and of Mubayi's result [26] already mentioned in the first remark at the end of the previous section is that if we require only that the three edges of at least one P_4 of every K_4 subgraph get trifferent in some 3-edge-coloring then the number of colorings we need is strictly less than the above mentioned trivial upper bound.

Definition 2. *Let $p(n)$ denote the minimum number of 3-edge-colorings of K_n needed if we want that for all 4 vertices the edge set of some P_4 subgraph on those 4 vertices gets totally multicolored in at least one of these colorings.*

Proposition 4.

$$p(n) \leq \frac{1}{2} c_{3,3} \log n + o(\log n).$$

Proof. Color the edges of K_n with $m := n^{1/2} e^{c\sqrt{\log n}} = n^{1/2+o(1)}$ colors with Mubayi's method so that at least 4 colors appear on every K_4 subgraph. This means that at least one P_4 of each K_4 is totally multicolored. Now encode each color with a ternary sequence

of length $t(m, 3, 3)$ so that any 3 sequences are trifferent in at least one coordinate. Now substituting every color in the previous coloring with a sequence of 3-colorings according to the ternary codes of colors, we will get a required sequence of 3-colorings, since any three edges that received three different colors in the first coloring must get three different colors in the 3-coloring that belongs to the position where the codewords of those 3 colors are trifferent.

Since $t(m, 3, 3) = c_{3,3} \log m + o(m) = c_{3,3} \frac{1}{2} \log n + o(\log n)$ the result follows. \square

Remark 3. A K_4 that is edge-colored with 4 colors contains at least 4 totally multicolored P_4 's out of the altogether 12 ones. It would be interesting to see whether one could apply Mubayi's coloring 3 times in such a way that all P_4 's get 3-colored in at least one of them. If this would be possible, then the argument in the proof of Proposition 4 would give that with $\frac{3}{2}c_{3,3} \log n + o(\log n)$ 3-edge colorings of K_n one can totally multicolor all its P_4 's thus providing a strictly better upper bound than the trivial $2c_{3,3} \log n + o(\log n)$. \diamond

In the following theorem we use a variation of an idea from [23] to give an explicit construction that provides another upper bound for $p(n)$. Because of the lack of our knowledge of the value of $c_{3,3}$, we do not know whether this upper bound is better or worse than the one provided by the previous proposition.

Theorem 2.

$$p(n) \leq 2 \lceil \log_2 n \rceil.$$

Proof. First we give $\lceil \log_2 n \rceil$ 4-colorings of the edges of K_n with the property that any 4-element subset of the vertex set contains a P_4 subgraph that is totally multicolored in at least one of them.

Label the n vertices with the binary form of the numbers $0, 1, \dots, n-1$ padded with 0's at the beginning if necessary to make all these binary sequences having the same length $t := \lceil \log_2 n \rceil$. This way each vertex is attached a different 0 – 1 sequence of length t . Let u, v be two distinct vertices and u_i, v_i be the i th coordinate in the binary sequences attached to u and v , respectively. We define a coloring of edge u, v for every $1 \leq i \leq t$ as follows. Color the edge $\{u, v\}$

- blue if $u_i = v_i = 0$,
- green if $u_i = v_i = 1$,
- red if $u_i \neq v_i$ and $\forall j < i : u_j = v_j$,
- yellow if $u_i \neq v_i$ and $\exists j < i : u_j \neq v_j$.

Consider any set of 4 vertices a, b, c, d . If there is any coordinate i where exactly two of a_i, b_i, c_i, d_i are 0 and two are 1, then we have two non-incident edges one of which is blue, while the other is green, and the remaining 4 edges of the K_4 with vertices a, b, c, d are all red and yellow. So those P_4 's (there are 4 of them), that contain the blue and green edges are totally multicolored.

If there is no such coordinate where exactly two of these four sequences are 0, then there are at least three such coordinates where one of the values is different from the other three. Without loss of generality we may assume that $a_i \neq b_i = c_i = d_i, b_j \neq a_j = c_j = d_j$,

and $c_h \neq a_h = b_h = d_h$, where we may also assume $i < j < h$ and that these i, j, h are the smallest coordinates where the above relations hold. Then the coloring rule gives that in the j th coloring the edge $\{a, b\}$ is yellow, $\{b, c\}$ is red, and $\{c, d\}$ is either blue or green, thus giving a totally multicolored P_4 . In fact, also the P_4 with edges $\{a, b\}, \{b, d\}, \{c, d\}$ is totally multicolored in the j th coloring, while two other P_4 's will be totally multicolored in the h th coloring, so altogether we can guarantee at least 4 different totally multicolored P_4 's on any given quadruple of vertices.

Since $t(4, 3, 3) = 2$ (just take the 2-length sequences 00, 01, 12, 22, any three of these are trifferent in at least one of the two coordinates), we can encode the 4 colors with two length ternary sequences that make any three of them trifferent in at least one coordinate. Thus replacing each of the 4-colorings in the previous construction with two appropriate 3-colorings (just identifying the colors blue and green in one and the colors yellow and red in the other), we obtain $2\lceil \log_2 n \rceil$ 3-edge colorings satisfying that for every 4 vertices some P_4 on them is totally multicolored by at least one of these colorings. In fact, we have seen, that at least 4 out of the 12 P_4 's on any given four vertices will be rainbow at some coloring. \square

Naturally, a similar remark to that of Remark 3. above could be formulated concerning the coloring used in the proof of Theorem 2.

It should be clear from the proof of Theorem 2 that if we label the vertices of K_n with t -length binary sequences so, that for any four of them there is a coordinate where exactly two of these four sequences contain a 1, then $p(n) \leq t$. (We could simply use the same coloring strategy but identifying colors red and yellow.) Alon, Körner, and Monti [1] investigated the maximum number of binary sequences of some given length with this property. Their results give, that if $\mu(n)$ is the minimum t for which n such t -length sequences can be given, then $\frac{1}{0.78} \log_2 n + o(\log n) \leq \mu(n) \leq \frac{3}{\log_2 \frac{8}{5}} \log_2 n + o(\log n) \approx 4.425 \log_2 n + o(\log n)$. In a similar manner one can obtain an upper bound on the minimum number of 3-colorings needed to make every P_4 totally multicolored in at least one of them if we use so-called (2, 2)-separated binary sequences that means a set of sequences with the property that for any two disjoint pair of them, there is a coordinate where both sequences of the first pair contain a 0 and both sequences of the other pair contain a 1. The problem of maximum cardinality of such a set of binary sequences of some fixed length is considered, for example, in [16] and [22].

An immediate consequence of Proposition 4, Theorem 2, and the discussion in the previous paragraph is the following. Since we do not know whether $c_{3,3}$ is smaller or larger than 4 (cf. the bounds quoted at the beginning of this section for $c_{3,3}$), and the value of $\mu(n)$ is also unknown, we do not know which expression achieves the minimum in the statement.

Corollary 1.

$$p(n) \leq \min\{2\lceil \log_2 n \rceil, \frac{1}{2}c_{3,3} \log n + o(\log n), \mu(n)\}.$$

\square

It is rather annoying that we do not even know whether $p(n) = \Omega(\log n)$ is true. Since $p(n) \geq \log_3 a(n, P_4)$ and $a(n, P_4) \geq f(n, 4, 3)$, this would immediately follow if we knew a lower bound $f(n, 4, 3) \geq n^\varepsilon$ for any constant $\varepsilon > 0$. This is, however, not known, in fact, the behavior of $f(n, 4, 3)$ is mentioned as a most annoying open problem in [13], where a $c\sqrt{n}$ lower bound is given for it, while the authors also mention that they could not even prove that $\frac{f(n, 4, 3)}{\log n}$ tends to infinity.

3.2 Bicoloring the triangle and Shannon capacity

It is rather trivial to see that if we want to two-color only certain subgraphs in several rounds then the problem of determining the minimum number of rounds is equivalent to some of those problems discussed in Section 2, in particular in determining the value $b(n, F)$. Indeed, if we can use a fixed number b of colors in several rounds and we want that each subgraph F gets at least two colors in at least one of the rounds, then the minimum number of rounds needed is exactly $\lceil \log_b b(n, F) \rceil$.

Propositions 1 and 2 show that in case of $F = K_{1,3}$ and $F = P_4$ we have $\log b(n, F) = \Theta(\log n)$. The situation is rather different, however, for $F = K_3$, because $b(n, K_3) = \min\{k : R(3; k) > n\}$ and the growth rate of $R(3; k)$ is at least exponential in k . This means that the minimum number of colorings (even with two colors) making every triangle bicolored in at least one of them is bounded from above by $O(\log \log n)$. On the other hand, it is not known, whether the growth rate of $R(3; k)$ is exponential or even faster in k . This is a famous open problem of Erdős, who, according to [8] (see Chapter 2.5) offered \$250 for determining the limit $\lim_{k \rightarrow \infty} R(3; k)^{1/k}$ and \$100 for deciding whether it is finite. (The best lower bound on this limit we know about is $(321)^{1/5} \approx 3.171765\dots$ due to Exoo [15].) This question is well-known to be equivalent to a question about Shannon capacity, cf. [2, 14, 28].

Definition 3. *Let G and H be two graphs. Their OR-product $G \otimes H$ is defined as follows.*

$$V(G \otimes H) = V(G) \times V(H)$$

and

$$E(G \otimes H) = \{(g_1, h_1), (g_2, h_2)\} : (g_1, g_2) \in E(G) \text{ or } (h_1, h_2) \in E(H)\}.$$

G^t denotes the t -fold OR-product of graph G with itself.

Definition 4. *The Shannon OR-capacity of graph G is defined as the always existing limit*

$$C_{\text{OR}}(G) = \lim_{t \rightarrow \infty} \sqrt[t]{\omega(G^t)},$$

where ω stands for the clique number.

Shannon capacity is most often defined by another graph product (the so called AND-product) and via independence numbers, see [31, 25], but the two notions are essentially

the same as the Shannon OR-capacity of a graph G we defined here is simply the more usual Shannon capacity (the “Shannon AND-capacity”) of the complementary graph \bar{G} . We use this less often used version because certain statements are more natural in this language and it is more convenient for our goals. (This approach also appears in several papers, cf. for example [17, 30].) It is well-known that $\omega(G) \leq C_{\text{OR}}(G) \leq \chi(G)$, where $\chi(G)$ is the chromatic number of G . In general, Shannon (OR)-capacity is very hard to determine, the only odd cycle longer than 3 for which it is known is C_5 , and the determination of $C_{\text{OR}}(C_5)$ was a major result of Lovász [25]. It was also a sort of breakthrough when Bohman and Holzman [6] proved with an ingenious construction that the Shannon OR-capacity of every odd cycle is strictly larger than its trivial lower bound 2, i.e., $C_{\text{OR}}(C_{2k+1}) > 2$ for every positive integer k . Erdős, McEliece, and Taylor [14] (and independently but later Alon and Orlitsky [2], cf. also [28]) showed, that the supremum of the possible Shannon OR-capacities a triangle-free graph can have is equal to $\lim_{k \rightarrow \infty} \sqrt[k]{R(3; k)}$. That is, Erdős’s question whether this limit is infinite is equivalent to the question, whether a triangle-free graph can have arbitrarily large Shannon OR-capacity. Based on a nice observation by Cropper, Gyárfás, and Lehel [9] we claim that to decide this, it would be enough to know the Shannon OR-capacity of the well-known Mycielski graphs [27] that we define below.

Definition 5. For any graph G on vertex set $\{v_1, \dots, v_n\}$ its Mycielskian $M(G)$ is defined as follows.

$$V(M(G)) = V(G) \cup \{u_i : v_i \in V(G)\} \cup \{z\},$$

$$E(M(G)) = E(G) \cup \{\{u_i, v_j\} : \{v_i, v_j\} \in E(G)\} \cup \{\{u_i, z\} : v_i \in V(G)\}.$$

By the k th Mycielski graph M_k we mean the result of the $(k-2)$ th iteration of the operation $M(\cdot)$ defined above when starting with $M_2 := K_2$. The vertex u_i we refer to as the twin vertex of v_i .

We index M_k as defined here, because as is well-known, the Mycielski construction increases the chromatic number by 1 (while not changing the clique number) and this way M_k denotes the k -chromatic Mycielski graph. Note that M_3 is just C_5 and M_4 is the graph also called Grötzsch graph.

Theorem 3.

$$\lim_{k \rightarrow \infty} \sqrt[k]{R(3; k)} = \lim_{r \rightarrow \infty} C_{\text{OR}}(M_r).$$

Proof. By the result of Erdős-McEliece-Taylor [14] and Alon-Orlitsky [2] already quoted above we know that

$$\lim_{k \rightarrow \infty} \sqrt[k]{R(3; k)} = \sup\{C_{\text{OR}}(G) : K_3 \not\subseteq G\}.$$

Cropper, Gyárfás, and Lehel [9] observed that every triangle-free graph is an induced subgraph of M_r if r is large enough. Since $C_{\text{OR}}(G)$ is obviously monotone in the sense that $F \subseteq G$ implies $C_{\text{OR}}(F) \leq C_{\text{OR}}(G)$ and the graphs M_r are triangle-free themselves,

we have that $\sup\{C_{\text{OR}}(G) : K_3 \not\subseteq G\} = \lim_{r \rightarrow \infty} C_{\text{OR}}(M_r)$ and thus the statement of the theorem. (Note that the existence of the limit follows from $M_{r-1} \subseteq M_r$ and thus $C_{\text{OR}}(M_r) \geq C_{\text{OR}}(M_{r-1})$ being true for every r .) \square

We do not know how to show that the Shannon OR-capacity of Mycielski graphs does or does not go to infinity. As a modest step forward we give a lower bound on the Shannon OR-capacity of the Grötzsch graph that is strictly larger than $C_{\text{OR}}(C_5) = \sqrt{5}$, i.e, the Shannon OR-capacity of the previous Mycielski graph.

Theorem 4.

$$C_{\text{OR}}(M_4) \geq \sqrt[4]{28}.$$

Proof. As $C_{\text{OR}}(G) \geq \sqrt[t]{\omega(G^t)}$ for any finite t , it is enough to present a clique of size 28 in the 4th OR-power of the Grötzsch graph $M_4 = M(C_5)$. This is what we do. We denote the vertices of C_5 by $0, 1, 2, 3, 4$, where every vertex is connected to the vertices labeled by neighboring integers, plus $\{0, 4\}$ is also an edge. The corresponding twin vertices (cf. Definition 5) are labeled $0', 1', 2', 3', 4'$, respectively, and the eleventh vertex of $M(C_5)$ is labeled z as in Definition 5.

We claim that the following 28 vertices form a clique in $[M(C_5)]^2$. Indeed, it is enough to observe, that the 25 vertices listed in the first five rows below are created from Shannon's well-known 5-element clique of C_5^2 in [31] given as $00, 12, 24, 31, 43$ by taking its second power and then "lifting" some vertices to their twins in such a way that at least one connection still remained between any two of our sequences. All sequences except five have a lifted element in the last two coordinates and for the five exceptional ones (see them put diagonally) both of the first two coordinates are lifted. Thus all these sequences will be connected with the three in the last row that are also connected to each other.

$$\begin{array}{ccccc} 0'0'00 & 120'0 & 2400' & 3100' & 430'0 \\ 001'2 & 1'2'12 & 241'2 & 3112' & 4312' \\ 0024' & 122'4 & 2'4'24 & 312'4 & 4324' \\ 0031' & 1231' & 243'1 & 3'1'31 & 433'1 \\ 004'3 & 1243' & 2443' & 314'3 & 4'3'43 \\ & 0'zzz & z0'zz & zzzz & \end{array}$$

Thus we have $\omega([M(C_5)]^2) \geq 28$ and the statement follows. \square

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