CONGRUENCES ON CLOSED SETS OF SELF-DUAL FUNCTIONS IN MANY-VALUED LOGICS AND ON CLOSED SETS OF LINEAR FUNCTIONS IN PRIME-VALIED LOGICS

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A congruence on a closed set *M* in the *k*-valued logic  $P_k$ , *k*≥2, is an equivalence relation on *M* which is compatible with the operations (permutation and identification of variables, addition of fictitious variables and substitution) of  $P_k$ .

In this paper we prove that the congurences on closed sets of self-dual functions of  $P_k$  are determined by congruences on closed sets of non-self-dual functions.

Moreover, we determine all congruences on the closed sets of linear functions (see [1] and [8]) in prime-valued logic.

1. Basic concepts

Let  $E_k$  denote the set  $\{0, 1, \dots, k-1\}$ , where  $k \ge 2$ . Let  $P_k^n$  denote the set of all functions  $f^n: E_k^n \rightarrow E_k$   $(n \ge 1)$  and put  $P_k = \bigcup_{n\ge 1} P_k^n$ . If there is no danger of confusion, the super-

script n of the function  $f^n$  is omitted.

The set of all function of  $P_k^1$  having exactly 1 values we denote by  $P_k^{[1]}$ .

The functions  $e_i^n \mathcal{GP}_k$   $(1 \le i \le n)$  defined by  $e_i^n(x_1, \ldots, x_n) = x_i$ are called projections. The *n*-ary constant function with value a is denoted by  $e_n^n$ .

The operations on  $P_k$  are  $\zeta$ ,  $\tau$ ,  $\Delta$ ,  $\nabla$ , \*, which are defined by

 $\begin{aligned} & (\zeta f) (x_1, \dots, x_n) = f(x_2, x_3, \dots, x_n, x_1) \\ & (\tau f) (x_1, \dots, x_n) = f(x_2, x_1, x_3, \dots, x_n) \\ & (\Delta f) (x_1, \dots, x_{n-1}) = f(x_1, x_1, x_2, \dots, x_{n-1}) \\ & (\nabla f) (x_1, \dots, x_{n+1}) = f(x_2, x_3, \dots, x_{n+1}) \end{aligned}$ 

 $(f^*g)(x_1, x_2, \dots, x_{n+m-1}) = f(g(x_1, \dots, x_m), x_{m+1}, \dots, x_{n+m-1}),$ where  $f^n$ ,  $g^m \in P_k$  (see [5] or [6]).

Superpositions over the set  $A\_P_k$  are functions obtained from A by using the operations  $\zeta, \tau, \Delta, \nabla, *$  finitely many times. The closure [A] of a set  $A \subseteq P_k$  is the set of all superpositions over A. A set A is said to be closed if [A]=A.

An equivalence relation  $\kappa$  on a closed set A is called a congruence on A iff  $f \sim g(\kappa)$ ,  $s \sim t(\kappa)$  imply  $f \ast s \sim g \ast t(\kappa)$  and  $\alpha f \sim \alpha g(\kappa)$  for all  $\alpha \in \{\zeta, \tau, \Delta, \nabla\}$  and for all  $f, g, s, t \in A$ . A.I.Mal'cev showed in |5| that every closed set  $A_{-}P_{k}$  has three trivial congruences  $\kappa_{0}$ ,  $\kappa_{\alpha}$  and  $\kappa_{1}$ :

 $f \sim g \ (\kappa_0) : \iff f = g \Lambda \{f, g\} \underline{c} A$   $f^n \sim g^m \ (\kappa_a) : \iff n = m \Lambda \{f, g\} \underline{c} A$   $f \sim g \ (\kappa_1) : \iff \{f, g\} \underline{c} A.$ 

Let K and K' be two congruences on A. We write  $K \leq K'$  iff  $f \sim g$  (K) implies  $f \sim g$  (K') for all f,  $g \in A$ .

For the other undefined notations we refer the reader to [1]-[8], particularly to [6].

For the proofs of our theorems we need the following lemma which is well known.

1.1 LEMMA ([2]). Let A be a closed set in  $P_k$  containing the projections. If  $\kappa$  is a congruence on A with  $\kappa \not \leq \kappa_a$ , then  $\kappa = \kappa_1$ .

## 2. Congruences on closed sets of self-dual functions of Pk

Let  $s(x)=x+1 \mod k$  and let S be the set of all functions of  $P_k$  preserving the relation  $\{(a, s(a)) \mid a \in E_k\}$ . The functions of S are called self-dual functions. If k is a prime number then S is a maximal closed set of  $P_k$  ([7]). In [3] for a maximal closed set of self-dual functions it was proved that a function  $f^n \mathcal{GP}_k$  is a function of *S* iff there exists a function  $F^{n-1}\mathcal{GP}_k$  with the property

$$f(\tilde{x}) = \sum_{i=0}^{k-1} j_i(x_1) \cdot s^i(F(s^{k-i}(x_2), \dots, x^{k-i}(x_n))) \mod k, \quad (1)$$

where  $s^{i}(x)$ :  $x+i \mod k$ ;  $j_{i}(x) = \begin{cases} 1 & if \ x & i \\ 0 & otherwise \end{cases}$  and

 $F(x_1, \ldots, x_{n-1}) = f(0, x_1, \ldots, x_{n-1})$ . The proof in [3] does not use the property that k is prime. Therefore (1) is right for every k. Because of this property we can define a bijective mapping  $\alpha$  of  $S(\_P_k)$  onto  $P'_k: \{f^n: E_k^n \to E_k\}$  as follows:

 $\alpha : f \rightarrow F.$ 

2.1 LEMMA. The mapping  $\alpha$  has the following properties:

(i) For the operations  $\hat{\zeta}, \hat{\tau}, \hat{\Delta}, \hat{\nabla}, \hat{*}$  defined by

 $(\hat{\zeta}f) (x_1, \dots, x_n) = f(x_1, x_3, x_4, \dots, x_n, x_2)$  $(\hat{\tau}f) (x_1, \dots, x_n) = f(x_1, x_3, x_2, x_4, \dots, x_n)$  $(\hat{\Delta}f) (x_1, \dots, x_{n-1}) = f(x_1, x_2, x_2, x_3, \dots, x_{n-1})$  $(\hat{\nabla}f) (x_1, \dots, x_{n+1}) = f(x_1, x_3, x_4, \dots, x_n)$  $(f^*g) (x_1, \dots, x_{n+m-2}) = f(x_1, g(x_1, x_2, \dots, x_m), x_{m+1}, \dots, x_{n+m-2})$  $is \alpha(\hat{\zeta}f) = \zeta F, \alpha(\hat{\tau}f) = \tau F, \alpha(\hat{\Delta}f) = \Delta F, \alpha(\hat{\nabla}f) = \nabla F \text{ and}$  $\alpha(f^*g) = F^*G, \text{ i.e. the algebra } \langle S; \hat{\zeta}, \hat{\tau}, \hat{\Delta}, \hat{\nabla}, \hat{\ast} \rangle \text{ is isomorphic} \\ \text{to the algebra } \langle P'_k; \zeta, \iota, \Delta, \nabla, \hat{\ast} \rangle.$ 

(ii) For every closed subset  $A(\neq \emptyset)$  of S is  $\alpha(A)$  a closed set,  $\alpha(A) \not\in S$  and  $A \in \alpha(A)$ .

PROOF. (i) is easy to check. Let A be a closed subset of S. Then by (i) we get that  $\alpha(A)$  is likewise a closed set. Assume  $\alpha(A) \subseteq S$ . Then

 $F(x_2,...,x_n) = s^i (F(s^{k-i}(x_2),...,s^{k-i}(x_n)))$ 

for  $i=0,1,\ldots,k-1$  and for every  $f^n \in A$  (see [3]). Thus by (1) we get that the variable  $x_1$  in every function  $f \in A$  is fictitious. however, this is not possible. Therefore is  $\alpha(A) \not \in S$ . Let  $f \in A$ . Then is  $\nabla f \in A$  and therefore  $\alpha(\nabla f) = f \in \alpha(A)$ , i.e.  $A \subset \alpha(A)$ .

2.2 THEOREM. Let A be a closed subset of  $S, \kappa$  a congruence on A and let  $\alpha(\kappa)$  defined by

 $F \sim G(\alpha(\kappa)) : \longleftrightarrow \alpha^{-1} F \sim \alpha^{-1} G(\kappa)$ 

an equivalence relation on  $\alpha(A)$ . Then

(i)  $\alpha(\kappa)$  is a congruence on  $\alpha(A)$  and

(ii)  $\alpha(\kappa)_{A}$  , i.e. the congruences on A we can get by restriction of the congruences on  $\alpha(A)$  to A.

PROOF. (i). Since  $\ltimes$  is a congruence on  $A \ltimes$  is also compatible with the operations  $\hat{\zeta}, \hat{\tau}, \hat{\Delta}, \hat{\nabla}, \hat{\ast}$ . Then by 2.1 (i) follows that  $\alpha(\ltimes)$  is a congruence on  $\alpha(A)$ . (ii) By 2.1 (ii) is  $A \subseteq \alpha(A)$  and therefore  $\alpha(\ltimes)_{/A}$  is a congruence on A. Let f and g be functions of A. If  $f^{\sim}g(\ltimes)$  then  $\nabla f_{\bullet} \nabla g(\ltimes)$  and by definition of  $\alpha(\ltimes)$  is  $\alpha(\nabla f) = f_{\bullet}g = \alpha(\nabla g)(\alpha(\ltimes)_{/A})$ , i.e.  $\underline{\subseteq} \alpha(\ltimes)_{/A}$ . If  $f_{\sim g}(\alpha(\ltimes)_{/A})$  then by definition of  $\alpha(\ltimes)$ we get that  $\alpha^{-1}f_{\sim} \alpha^{-1}g(\ltimes)$ . Since f and g are functions of s is  $\alpha^{-1}f = \nabla f$  and  $\alpha^{-1}g = \nabla g$ . Therefore is  $\nabla f_{\sim} \nabla g(\ltimes)$  and  $\Delta(\nabla f) = f_{\sim}g = \Delta(\nabla g)(\ltimes)$ , i.e.  $\alpha(\ltimes)_{/A} \subseteq \ltimes$ . Thus  $\alpha(\ltimes)_{/A} =$ .

## 3. Congruences on some closed subsets of $[P_k^I]$

In this section we prove a theorem which we need for the determination of the congruences on the closed subsets of linear functions.

Let  $C \subseteq P_k^{[1]}$ , G a subgroup os  $\langle P_k^{[k]}; * \rangle$ , where the functions of G preserve the set C, U a normal subroup of the group G and let  $\mu$  be an equivalence relation on C which is preserved by the functions of G.

It is easy to check that the equivalence relation  $U^{\mu}$ , on [GUC] defined by

$$f^{n} - g^{m}(\kappa^{U, \mu}) : \longleftrightarrow n = m\Lambda(\exists i \exists f', g' \in GUC: f(\tilde{x}) = f'(x_{i})\Lambda$$
$$g(\tilde{x}) = g'(x_{i})\Lambda(f' * U = g' * UVf' - g'(\mu)))$$

is a congruence on [GUC].

3.1 THEOREM. Let G be a subgroup of  $\langle P_k^{[k]}; * \rangle$ ,  $C \subseteq P_k^{[1]}$  and G  $\subseteq Pol C$ . Then exactly on [GUC] there exist the congruences  $\kappa_0, \kappa_a, \kappa_1$  and congruences of the type  $\kappa^{U,\mu}$ , where U is a normal subgroup of G and  $\mu$  is an equivalence relation on G which is preserved by the functions of G.

PROOF. Let  $\kappa$  be a congruence on [GUC] and  $\kappa \neq \kappa_1$ . Then by 1.1 is  $\kappa \subseteq \kappa_a$ . We have to ditinguish the following cases: Case 1: There exist  $\kappa$ -congruent functions  $f^n$  and  $g^n$  with  $\Delta^{n-1}f \in G$  and  $\Delta^{n-1}g \in C$ . Then is  $\Delta^{n-1}f \sim \Delta^{n-1}g(\kappa)$ . Thus  $e_1^1 \sim \Delta^{n-1}g =: e_a(\kappa)$ ,  $a \in E_k$ . By this we have for every function  $h^m \in [GUC]$ :  $e_1^1 * h = h \sim e_a^m = e_a * h(\kappa)$ , *i.e.*  $\kappa = \kappa_a$ . Case 2: There exist  $\kappa$ -congruent functions  $f^n$  and  $g^n$  with  $f(x_1, \ldots, x_n) = f'(x_i), g(x_1, \ldots, x_n) = g'(x_j), \{f', g'\} \subseteq G$ and  $i \neq j$ .

Without loss of generality we can assume that i=1 and j=2. The inverse functions of f' and g' we denote by f'' and g'', respectively. Then we have

$$\begin{split} f(f''(x_1),g''(x_2),x_2,\ldots,x_2) &= e_1^2(x_1,x_2) \\ &\sim g(f''(x_1),g''(x_2),x_2,\ldots,x_2) &= e_2^2(x_1,x_2) \quad (\texttt{K}). \text{ Therefore is} \\ &e_1^2(s(\widetilde{x}),t(\widetilde{x})) &= s(\widetilde{x}) \sim t(\widetilde{x}) &= e_2^2(s(\widetilde{x}),t(\widetilde{x})) \quad (\texttt{K}) \text{ for every } s \\ \text{and } t \text{ of } [GUC]^m, m \geq 1, \ i.e. \ \texttt{K} &= \texttt{K}_a. \end{split}$$

Case 3: Two *n*-ary functions f and g are  $\kappa$ -congruent if and only if either  $\{f,g\} \subseteq [C]$  or there exist an i and  $f',g' \in G$  with  $f(x_1,\ldots,x_n) = f'(x_i), g(x_1,\ldots,x_n) = g'(x_i).$ 

In this case the congruence  $\kappa$  is exactly determined by  $\kappa_{/G}$  and  $\kappa_{IC}$ .

As we know, the congruence on a group G are determined by a normal group U of G and  $f \sim g($  ) iff  $f^*U = g^*U$  for all f,  $g \notin G$ .

Obviously, C is an equivalence relation on C which is preserved by every function of G. Thus = U, C.

## 4. Congruences on closed sets of linear functions in prime-valued logics

Let p be a fixed prime number. L denote the set of all linear functions over  $\langle E_p; +, \cdot \mod p \rangle$  in  $P_p$ , i.e.

$$L:=\bigcup_{n>1} \{f^n \mathcal{GP}_p | \exists a_i: f(\tilde{x}) = a_0 + \sum_{i=1}^n a_i \cdot x_i \mod p\}.$$

In |1| it was proved that L has exactly the following closed subsets:

$$LAS = \bigcup_{n \ge 1} \{f^{n} \in L \mid a_{1} + a_{2} + \dots + a_{n} = 1 \mod p\},$$

$$LAPol(a) = \bigcup_{n \ge 1} \{f^{n} \in P_{p} \mid \exists a_{i} : f(\tilde{x}) = a + \sum_{i=1}^{n} a_{i} \cdot (x_{1} - a)\},$$

$$a \in E_{n},$$

LASAPol(0) and closed subsets A with  $A \leq [L^{1}]$ . If  $A \leq [L^{[1]}]$ , then it is easy to see that the closed set A has only congruences of the type  $\kappa^{\mu}$  and of the type  $\kappa^{\mu}_{a}$  defined by

$$f^{n} \sim g^{m}(\kappa^{\mu}): \longleftrightarrow \Delta^{n-1} f \sim \Delta^{m-1} g(\mu) \quad \text{and}$$
  
$$f^{n} \sim g^{m}(\kappa^{\mu}_{a}): \longleftrightarrow n = m \Lambda \Delta^{n-1} f \sim \Delta^{n-1} g(\mu),$$

where  $\mu$  is an any equivalence relation on  $A^{1}$ . If  $A \notin [L^{[1]}]$  and  $A \subseteq [L^{1}]$  then the congruences on A follow by theorem 3.1.

We denote by  $\kappa_c$  an equivalence relation defined by

$$f^{n} \sim g^{m}(\kappa_{c}) : \longleftrightarrow n = m\Lambda(\exists a: f(\tilde{x}) = a + g(\tilde{x}) \mod p).$$

Obviously,  $\kappa_c$  is a congruence on *L*. We will show that  $\kappa_c$  is the only nontrivial congruence on  $A \subseteq L$  for  $A \notin [L^1]$ .

- 4.1 THEOREM ([4]).  $\kappa_0$ ,  $\kappa_c$ ,  $\kappa_a$  and  $\kappa_1$  are the only congruences on L.
- 4.2 THEOREM.  $\kappa_0$ ,  $\kappa_a$  and  $\kappa_1$  are the only congruences on  $L\Omega Pol(a)$  for every  $a \in E_p$ .

PROOF. Clearly, the closed sets  $L\cap Pol(a)$ ,  $a \in E_p$ , are mutually isomorphic. Therefore we can assume that a=0. Let  $\kappa$  be a congruence on  $L\cap Pol(0)$ . The following two cases are possible:

Case 1:  $\kappa \not \leq \kappa_a$ . By 1.1 is  $\kappa = \kappa_1$ . Case 2:  $\kappa_0 \not \subset \kappa \not \subseteq -\kappa_a$ .

Then there exist  $\kappa$  -congruent functions  $f^n$ ,  $g^n$  and an n-tuple  $\tilde{a}=(a_1,\ldots,a_n)$  with  $f(\tilde{a})\neq g(\tilde{a})$ . Therefore is  $f(a_1\cdot x,a_2\cdot x,\ldots,a_n\cdot x)=:a\cdot x \cdot b\cdot x :=g(a_1\cdot x,a_2\cdot x,\ldots,a_n\cdot x)$  ( $\kappa$ ), where  $a\neq b$ .

The functions  $h(x,y)=x-y \mod p$  and  $t(x)=(a-b)^{-1}\cdot x$  belong to  $L\cap Pol(0)$ . Thus we get  $h(ax,ax)=c_0\sim h(ax,bx):=h'(x)$  () and  $c_0^1=t*c_0^1\sim t*h'=c_1^1$  (K). This implies that  $c_0^1*r^m=c_0^m\sim r=c_1^1*r$  (K) for every  $r^m \in L\cap Pol(0)$ ,  $m \ge 1$ . Therefore  $K = K_a$ .

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4.3 THEOREM.  $\kappa_0$ ,  $\kappa_c$ ,  $\kappa_a$  and  $\kappa_1$  are the only congruences on LOS.

**PROOF.** Obviously,  $\alpha(LAS)=L$ . Therefore, using 2.2 and 4.1 we have the theorem.

**4.4** THEOREM.  $\kappa_0$ ,  $\kappa_a$  and  $\kappa_1$  are the only congruences on LOSOPOl(0).

**PROOF.** The theorem follows from  $\alpha(LASAPol(0)) = LAPol(0)$ , 2.2 and 4.2.

We also remark that the structure of the congruences becomes more complicated, if k is not a prime number. If k is square-free, then follows by [8] and by [4] (theorem 3.6) that every closed subset of L has only a finite number of congruences. But, if k is not square-free then there exist closed subsets of L with an infinite number of congruences. Finally we give an example for a closed subset with a such a property.

Let Z:=  $\bigcup_{n \ge 1} \{f^n \in P_4 | \exists a_i \in \{0, 2\} : f(\tilde{x}) = \sum_{i=1}^n a_i \cdot x_i \mod p\}$ , let r(f)

be the number of the non-fictitious variables of the function f.

Further let  $\chi_i$  be an equivalence relation defined by

$$n^{n} \sim g^{m}(\chi_{i}): \iff f = g \vee (n = m \Lambda r(f) \leq i \Lambda r(g) \leq i),$$

 $i=1,2,\ldots$ . If is easy to prove that  $\chi_i$  for all  $i\geq 1$  is a congruence on Z of  $P_A$ .

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