

# A THEOREM OF FINITE SETS

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## § 1. Introduction

Let  $A_1, \dots, A_n$  be a system of different subsets of a finite set  $H$ , where  $|H| = h$  and  $|A_i| = l$  ( $1 \leq i \leq n$ ) ( $|A|$  denotes the number of elements of  $A$ ). We ask for a system  $A_1, \dots, A_n$  (for given  $h, l, n$ ) for which the number of sets  $B$  satisfying  $|B| = l - 1$  and  $B \subset A_i$  for some  $i$  is minimum. The first lower estimation for this minimum is given by SPERNER ([1], Hilfssatz).

His estimation is  $\frac{n \cdot l}{h - l + 1}$ . This depends on  $h$ . However, if  $n = \binom{N}{l}$ , it is expected that the minimizing system is the system of all  $l$ -tuples chosen from a subset of  $N$  elements of  $H$ . In this case the number of  $B$ 's is  $\binom{N}{l - 1}$

which does not depend on  $h$ . A. HAJNAL proved this statement in the case of  $l = 3$  (unpublished). In this paper I prove for all cases that this is, indeed, the minimum, and find the (more complicated) minimum also for arbitrary  $n$ . The theorem is probably useful in proofs by induction over the maximal number of elements of the subsets in a system, as was SPERNER's lemma in his paper [1].

KLEITMAN told me in Tihany (Hungary) that he thought I could solve the following problem of ERDŐS by the aid of the above theorem and the "marriage problem": Let  $A_1, \dots, A_n$  be subsets of  $H$ , where  $|H| = 2h$  and  $|A_i| = h$ . For what  $n$ 's is it always possible to construct a system  $B_1, \dots, B_n$  with the properties  $B_i \subset A_i$ ,  $|B_i| = h - 1$  ( $1 \leq i \leq n$ ). § 3 contains the solution of this problem in a more general form.

## § 2. The main result

Before the exact formulation of the theorem we need the following simple but interesting

LEMMA 1. *If  $n$  and  $l$  are natural numbers, we can write the number  $n$  uniquely in the form*

$$(1) \quad n = \binom{a_l(n, l)}{l} + \binom{a_{l-1}(n, l)}{l-1} + \dots + \binom{a_{t(n, l)}(n, l)}{t(n, l)},$$

where  $t(n, l) \geq 1$ ,  $a_l > a_{l-1} > \dots > a_{t(n, l)}$  are natural numbers and  $a_i(n, l) \geq i$  ( $i = t(n, l), t(n, l) + 1, \dots, l$ ).

PROOF. The existence of form (1) is proved by induction over  $l$ . For  $l = 1$  the statement is trivial. Assume that for  $l = k - 1$  it is true also and prove for  $l = k$ . Let  $a_k$  be the maximal integer satisfying the inequality

$\binom{a_k}{k} \leq n$ . If here equality holds, we are ready. If it does not, using the induction hypothesis we have for the number  $n - \binom{a_k}{k}$  the following expression:

$$(2) \quad n - \binom{a_k}{k} = \binom{a_{k-1}}{k-1} + \dots + \binom{a_t}{t},$$

where  $t \geq 1$ ,  $a_{k-1} > \dots > a_t$ ,  $a_i \geq i$  ( $i = t, t+1, \dots, k-1$ ). (2) gives an expression for  $n$ , we have to verify only  $a_k > a_{k-1}$  and  $a_k \geq k$ . If  $a_k \leq a_{k-1}$  held, then

$$n \geq \binom{a_k}{k} + \binom{a_{k-1}}{k-1} \geq \binom{a_k}{k} + \binom{a_k}{k-1} = \binom{a_k+1}{k}$$

would hold also, which contradicts choosing of  $a_k$ . On the other hand,  $a_k \geq k$  follows from  $a_k > a_{k-1}$  and  $a_{k-1} \geq k-1$ .

The unicity of Form (1) is proved also by induction over  $l$ . For  $l = 1$  the statement is trivial. Assume that for  $l = k - 1$  it is also true and prove for  $l = k$ . If, on the contrary, there exist two forms:

$$(3) \quad n = \binom{a_k}{k} + \binom{a_{k-1}}{k-1} + \dots + \binom{a_t}{t} = \binom{a'_k}{k} + \binom{a'_{k-1}}{k-1} + \dots + \binom{a'_r}{r},$$

we may separate two different cases. If  $a_k = a'_k$ , we can obtain two different forms of  $n - \binom{a_k}{k}$ , which contradict our induction hypothesis. If  $a_k < a'_k$ , the contradiction follows from

$$\begin{aligned} n &\leq \binom{a_k}{k} + \binom{a_k-1}{k-1} + \dots + \binom{a_k-k+1}{1} = \binom{a_k+1}{k} - 1 < \binom{a_k+1}{k} \leq \\ &\leq \binom{a'_k}{k} \leq \binom{a'_k}{k} + \dots + \binom{a'_r}{r}. \end{aligned}$$

Thus we proved the lemma.

In the future we will use the following two notations:

$$E_l(n) = \binom{a_l(n, l) - 1}{l-1} + \binom{a_{l-1}(n, l) - 1}{l-2} + \dots + \binom{a_{l(n, l)}(n, l) - 1}{l(n, l) - 1}$$

and

$$F_l(n) = \binom{a_l(n, l)}{l-1} + \binom{a_{l-1}(n, l)}{l-2} + \dots + \binom{a_{l(n, l)}(n, l)}{l(n, l) - 1}.$$

These numbers are uniquely determined by Lemma 1.

Let us consider now the problem. Let  $H$  be a finite set with  $h$  elements, and

$$\mathcal{A} = \{A_1, \dots, A_n\}$$

a system of different subsets of  $H$ , where the number of elements of  $A_i$  is

$$|A_i| = l \quad (1 \leq i \leq n).$$

Obviously,  $l$  is a fixed integer between 1 and  $h$ . Let  $c(\mathcal{A})$  denote the following system

$$c(\mathcal{A}) = \{B : |B| = l - 1 \text{ and } B \subset A_j \text{ for at least one } j\}.$$

The problem is to determine the minimum of  $|c(\mathcal{A})|$ , if  $h$ ,  $n$  and  $l$  are given. Theorem 1 gives the exact solution of this problem.

**THEOREM 1.** *Let  $h$ ,  $n$  and  $l$  be given integers with the properties*

$$h \geq 1, \quad 1 \leq l \leq h \quad \text{and} \quad 1 \leq n \leq \binom{h}{l}.$$

If  $H$  is a set of  $h$  elements, and

$$\mathcal{A} = \{A_1, \dots, A_n\}, \quad |A_i| = l \quad (i = 1, \dots, n)$$

a system of different subsets of  $H$ , then

$$\min |c(\mathcal{A})| = F_l(n),$$

where the minimum runs over all such systems  $\mathcal{A}$ .

**REMARK.** It is interesting, that  $\min |c(\mathcal{A})|$  does not depend on  $h$ . For example, SPERNER's estimation [1]:

$$c(\mathcal{A}) \geq \frac{n \cdot l}{h - l + 1}$$

depends on  $h$ .

Before the proof we shall give another theorem. We will prove them together.

**THEOREM 2.** *Let  $h$ ,  $n$  and  $l$  be given integers with the properties*

$$h \geq 1, \quad 1 \leq l \leq h \quad \text{and} \quad \binom{h}{l} \leq n \leq 2 \binom{h}{l}.$$

Further  $G$  and  $H$  are disjoint sets of  $h$  elements. If

$$\mathcal{A} = \{A_1, \dots, A_n\}$$

is a system of  $A_i$ 's, where

$$A_i \subset G \quad \text{or} \quad A_i \subset H \quad (1 \leq i \leq n)$$

and

$$|A_i| = l \quad (1 \leq i \leq n),$$

then

$$\min |c(\mathcal{A})| = \binom{h}{l-1} + F_l \left( n - \binom{h}{l} \right).$$

**PROOF. 1.** First we construct the minimizing system of Theorem 1. Denote this system by  $\mathcal{M}(h, n, l)$ . Obviously, it is sufficient to construct the system  $\mathcal{M}(a_i^*(n), n, l)$ , where  $a_i^*(n)$  is the least integer satisfying

$$\binom{a_i^*(n)}{l} \geq n.$$

The construction will be carried out by induction over  $l$ . If  $l = 1$ ,  $a_1^*(n) = n$  and  $\mathcal{M}(a_1^*(n), n, 1)$  consists of all the sets of one element. Assume we constructed already the system  $\mathcal{M}(a_{l-1}^*(n), n, l-1)$  for all  $n$ . Construct now  $\mathcal{M}(a_l^*(n), n, l)$ . If  $n = \binom{a_l(n, l)}{l}$ , then the minimizing system consists of all the subsets having  $l$  elements. If  $n > \binom{a_l(n, l)}{l}$ , let  $H$  be a set of  $a_l^*(n) = a_l(n, l) + 1$  elements, and  $e$  an element of  $H$ . Since  $a_l > a_{l-1}$ , we can construct the system  $\mathcal{M}\left(a_l(n, l), n - \binom{a_l(n, l)}{l}, l-1\right)$  on  $H - \{e\}$  by the induction hypothesis. Define the system  $\mathcal{N}$  in the following manner:

$$\mathcal{N} = \left\{ N \cup \{e\} : N \in \mathcal{M}\left(a_l(n, l), n - \binom{a_l(n, l)}{l}, l-1\right) \right\}.$$

If  $\mathcal{S}$  denotes the system of all subsets of  $H - \{e\}$ , having  $l$  elements, then  $\mathcal{S}$  and  $\mathcal{N}$  form together the system  $\mathcal{M}(a_l^*(n), n, l)$ . Indeed, the number of sets is  $\binom{a_l(n, l)}{l} + n - \binom{a_l(n, l)}{l} = n$  and we have only to verify

$$(4) \quad |c(\mathcal{M}(a_l^*(n), n, l))| = \binom{a_l(n, l)}{l-1} + \dots + \binom{a_{l(n, l)}(n, l)}{l(n, l) - 1} = F_l(n).$$

However, it is easy to see, that

$$|c(\mathcal{M}(a_l^*(n), n, l))| = \binom{a_l(n, l)}{l-1} + \left| c\left(\mathcal{M}\left(a_l(n, l), n - \binom{a_l(n, l)}{l}, l-1\right)\right) \right|$$

and by the induction hypothesis

$$\left| c\left(\mathcal{M}\left(a_l(n, l), n - \binom{a_l(n, l)}{l}, l-1\right)\right) \right| = \binom{a_{l-1}(n, l)}{l-2} + \dots + \binom{a_{l(n, l)}(n, l)}{l(n, l) - 1},$$

which proves (4).

2. The minimizing system of Theorem 2 consists of a complete system in  $G$ , and  $\mathcal{M}\left(h, n - \binom{h}{l}, l\right)$  in  $H$ .

3. In the previous two points we showed that in the case of Theorem 1

$$\min |c(\mathcal{A})| \leq F_l(n),$$

and in the case of Theorem 2

$$\min |c(\mathcal{A})| \leq \binom{h}{l-1} + F_l \left( n - \binom{h}{l} \right).$$

Thus, it is sufficient to verify

$$(5) \quad |c(\mathcal{A})| \geq F_l(n)$$

and

$$(6) \quad |c(\mathcal{A})| \geq \binom{h}{l-1} + F_l \left( n - \binom{h}{l} \right),$$

respectively. These statements will be proved by induction over  $l$ . If  $l = 1$ , both statements are trivial. Assume we have proved for all numbers  $< l$  and prove for  $l$ .

4. First we prove the inequality

$$(7) \quad F_l(n) \leq F_l(n_1) + F_{l-1}(n_2),$$

if

$$(8) \quad n = n_1 + n_2, \quad n_1 \geq 0, \quad n_2 \geq 0$$

are integers, and

$$(9) \quad n_2 \leq E_l(n).$$

The statement will be proved for fixed  $l$  and for every  $n, n_1, n_2$  using the induction hypothesis for  $l - 1$ . For the sake of simplicity we use the following notations:

$$\begin{aligned} t = t(n, l) & \quad a_i = a_i(n, l) & \quad (t \leq i \leq l) & \quad a_i^* = a_i^*(n), \\ r = t(n_1, l) & \quad b_i = a_i(n_1, l) & \quad (r \leq i \leq l) & \quad b_i^* = a_i^*(n_1), \\ s = t(n_2, l - 1) & \quad c_i = a_i(n_2, l - 1) & \quad (s \leq i \leq l - 1) & \quad c_{i-1}^* = a_{i-1}^*(n_2). \end{aligned}$$

It follows from (8) and (9) that

$$(10) \quad n_1 \geq n - E_l(n) = \binom{a_l - 1}{l} + \dots + \binom{a_l - 1}{t}.$$

Because of (10)

$$(11) \quad b_l \geq a_l - 1$$

must hold, since in the contrary case it would be

$$n_1 \leq \binom{a_l - 2}{l} + \binom{a_l - 3}{l - 1} + \dots + \binom{a_l - (l + 1)}{1} = \binom{a_l - 1}{l} - 1,$$

what contradicts (10). On the other hand

$$(12) \quad a_l \geq b_l$$

because of (8). Applying (11) and (12) we can distinguish two different cases:  
 (a)  $b_l = a_l$  and (b)  $b_l = a_l - 1$ .

(a) In this case (7) has the form

$$\begin{aligned} & \binom{a_l}{l-1} + \binom{a_{l-1}}{l-2} + \dots + \binom{a_t}{t-1} \leq \\ & \leq \binom{a_l}{l-1} + \binom{b_{l-1}}{l-2} + \dots + \binom{b_r}{r-1} + F_{l-1}(n_2). \end{aligned}$$

Decreasing both sides by  $\binom{a_l}{l-1}$  we have

$$(13) \quad F_{l-1}\left(n - \binom{a_l}{l}\right) \leq F_{l-1}\left(n_1 - \binom{a_l}{l}\right) + F_{l-1}(n_2).$$

Let  $H_1$  and  $H_2$  be disjoint sets. Construct the system  $\mathcal{M}\left(b_{l-1} + 1, n_1 - \binom{a_l}{l}, l-1\right)$  on  $H_1$  and the system  $\mathcal{M}(c_{l-1}^*, n_2, l-1)$  on  $H_2$ . In this manner we obtain a system  $\mathcal{N}$  on  $H_1 \cup H_2$ . Applying the induction hypothesis (Point 3. (5)) for  $\mathcal{N}$  and  $l-1$  we have

$$(14) \quad \begin{aligned} & F_{l-1}\left(n - \binom{a_l}{l}\right) \leq |c(\mathcal{N})| = \\ & = \left| c\left(\mathcal{M}\left(b_{l-1} + 1, n_1 - \binom{a_l}{l}, l-1\right)\right) \right| + |c(\mathcal{M}(c_{l-1}^*, n_2, l-1))|. \end{aligned}$$

However, we know (Point 1. (4)) that

$$(15) \quad \left| c\left(\mathcal{M}\left(b_{l-1} + 1, n_1 - \binom{a_l}{l}, l-1\right)\right) \right| = F_{l-1}\left(n_1 - \binom{a_l}{l}\right)$$

and

$$(16) \quad |c(\mathcal{M}(c_{l-1}^*, n_2, l-1))| = F_{l-1}(n_2).$$

Finally, (13) follows from (14), (15), and (16).

(b)  $b_l = a_l - 1$ . We separate this case into two subcases:

$$(ba) \quad n_2 \geq \binom{a_l - 1}{l-1}, \quad (bb) \quad n_2 < \binom{a_l - 1}{l-1}.$$

(ba) In this case (7) has the form

$$\binom{a_l}{l-1} + \binom{a_{l-1}}{l-2} + \dots + \binom{a_t}{t-1} \leq \binom{a_l-1}{l-1} + \binom{b_{l-1}}{l-2} + \dots + \binom{b_r}{r-1} + \binom{a_l-1}{l-2} + \binom{c_{l-2}}{l-3} + \dots + \binom{c_s}{s-1},$$

since  $c_{l-1} = a_l - 1$ , because of (9) and the supposition (ba). Decreasing both sides by  $\binom{a_l}{l-1} = \binom{a_l-1}{l-1} + \binom{a_l-1}{l-2}$  we have

$$(17) \quad F_{l-1}\left(n - \binom{a_l}{l}\right) \leq F_{l-1}\left(n_1 - \binom{a_l-1}{l}\right) + F_{l-2}\left(n_2 - \binom{a_l-1}{l-1}\right).$$

We can prove (17) by using of the induction hypothesis if

$$(18) \quad n_2 - \binom{a_l-1}{l-1} \leq E_{l-1}\left(n - \binom{a_l}{l}\right)$$

holds. However (9) gives

$$(19) \quad \binom{a_l-1}{l-1} + \binom{c_{l-2}}{l-2} + \dots + \binom{c_s}{s} \leq \binom{a_l-1}{l-1} + \binom{a_{l-1}-1}{l-2} + \dots + \binom{a_t-1}{t-1}.$$

Decreasing both sides by  $\binom{a_l-1}{l-1}$  we obtain

$$(20) \quad \binom{c_{l-2}}{l-2} + \dots + \binom{c_s}{s} \leq \binom{a_{l-1}-1}{l-2} + \dots + \binom{a_t-1}{t-1}$$

and (20) is equivalent to (18).

(bb) In this case (7) has the form

$$\binom{a_l}{l-1} + \binom{a_{l-1}}{l-2} + \dots + \binom{a_t}{t-1} \leq \binom{a_l-1}{l-1} + \binom{b_{l-1}}{l-2} + \dots + \binom{b_r}{r-1} + F_{l-1}(n_2).$$

Decreasing both sides by  $\binom{a_l-1}{l-1}$  we have

$$(21) \quad \binom{a_l-1}{l-2} + F_{l-1}\left(n - \binom{a_l}{l}\right) \leq F_{l-1}\left(n_1 - \binom{a_l-1}{l}\right) + F_{l-1}(n_2).$$

Let  $G$  and  $H$  be two disjoint sets of  $a_l - 1$  elements. Construct the system  $\mathcal{M}\left(a_l - 1, n_1 - \binom{a_l - 1}{l}, l - 1\right)$ . We can construct it if  $n_1 - \binom{a_l - 1}{l} \leq \binom{a_l - 1}{l - 1}$ . But this follows from  $a_l - 1 = b_l > b_{l-1}$ , since

$$n_1 - \binom{a_l - 1}{l} = \binom{b_{l-1}}{l - 1} + \dots + \binom{b_r}{r}.$$

Construct further the system  $\mathcal{M}(c_{l-1}^*, n_2, l - 1)$  on  $H$ . The possibility of this construction follows from the assumption (bb). In this manner we obtain a system  $\mathcal{N}$  on  $G \cup H$ . Applying the induction hypothesis (Point 3. (6)) for  $\mathcal{N}$  and  $l - 1$  we have

$$\begin{aligned} \binom{a_l - 1}{l - 2} + F_{l-1}\left(n - \binom{a_l}{l}\right) &\leq \left| c\left(\mathcal{M}\left(a_l - 1, n_1 - \binom{a_l - 1}{l}, l - 1\right)\right) \right| + \\ (22) \quad &+ |c(\mathcal{M}(c_{l-1}^*, n_2, l - 1))|. \end{aligned}$$

However, we know (Point 1. (4)) that

$$(23) \quad \left| c\left(\mathcal{M}\left(a_l - 1, n_1 - \binom{a_l - 1}{l}, l - 1\right)\right) \right| = F_{l-1}\left(n_1 - \binom{a_l - 1}{l}\right)$$

and

$$(24) \quad |c(\mathcal{M}(c_{l-1}^*, n_2, l - 1))| = F_{l-1}(n_2),$$

further, (21) follows from (22), (23) and (24). Thus we proved the inequality for  $l$ .

5. However, we need (7) under the condition

$$(25) \quad n_2 \leq \frac{n \cdot l}{a_l^*}$$

instead of (9). Thus we are going now to prove the inequality

$$(26) \quad \frac{n \cdot l}{a_l^*} \leq E_l(n).$$

We prove (26) by induction over  $l$ , but we should like to mention that the proof of (26) is independent from the whole proof of the theorems. For  $l = 1$  the statement is trivial. Assume we proved it for the integers  $< l$ , and prove for  $l$ . If  $n = \binom{a_l}{l}$  then  $a_l^* = a_l$  and  $E_l(n) = \binom{a_l - 1}{l - 1}$ , thus (26) holds with equality. We may assume  $a_l^* = a_l + 1$ . Obviously

$$(27) \quad \frac{l}{a_l + 1} \binom{a_l}{l} \leq \binom{a_l - 1}{l - 1}$$

and by the induction hypothesis

$$(28) \quad \frac{l-1}{a_{l-1}+1} \left[ \binom{a_{l-1}}{l-1} + \dots + \binom{a_r}{r} \right] \leq \binom{a_{l-1}-1}{l-2} + \dots + \binom{a_r-1}{r-1}.$$

If  $\frac{l}{a_l+1} \leq \frac{l-1}{a_{l-1}+1}$ , summarizing (27) and (28) we obtain (26). In the contrary case

$$(29) \quad \frac{l}{a_l+1} > \frac{l-1}{a_l}$$

holds because of  $a_l \geq a_{l-1} + 1$ . Let us set out from the identity

$$\binom{a_l}{l-1} \left( \frac{l}{a_l+1} - \frac{l-1}{a_l} \right) = \binom{a_l-1}{l-1} - \frac{l}{a_l+1} \binom{a_l}{l}.$$

The expression in the bracket is positive because of (29), thus we can write

$$\left[ \binom{a_{l-1}}{l-1} + \binom{a_{l-2}}{l-2} + \dots + \binom{a_r}{r} \right] \left( \frac{l}{a_l+1} - \frac{l-1}{a_l} \right) < \binom{a_l-1}{l-1} - \frac{l}{a_l+1} \binom{a_l}{l},$$

since  $a_l > a_{l-1}$ . Write  $\frac{l-1}{a_{l-1}+1}$  instead of  $\frac{l-1}{a_l}$ , and reorder the inequality

$$\begin{aligned} \frac{l}{a_l+1} \left[ \binom{a_l}{l} + \binom{a_{l-1}}{l-1} + \dots + \binom{a_r}{r} \right] &< \binom{a_l-1}{l-1} + \\ &+ \frac{l-1}{a_{l-1}+1} \left[ \binom{a_{l-1}}{l-1} + \dots + \binom{a_r}{r} \right]. \end{aligned}$$

Finally, from the above inequality (26) follows by (28).

6. Now let us prove statement (5) for  $l$  by induction over  $h$  if  $h = l$  is trivial. Assume we have proved (5) for all sets  $|H| < h$ , and prove for  $h$ .

There exists an element  $e$  of  $H$ , contained by at most  $\frac{n \cdot l}{h}$  sets  $A_i$ . We

define the following systems:

$$\mathcal{B} = \{A : A \in \mathcal{A}, e \notin A\}$$

and

$$\mathcal{C} = \{A - \{e\} : A \in \mathcal{A}, e \in A\}$$

where

$$(30) \quad n_2 = |\mathcal{C}| \leq \frac{n \cdot l}{h} \leq \frac{n \cdot l}{a_l^*(n)}.$$

Naturally,

$$c(\mathcal{B}) \subset c(\mathcal{A})$$

and

$$c(\mathcal{C}) \cup e \subset c(\mathcal{A}),$$

where  $\mathcal{D}(\cup)a$  denotes in general the system  $\{D \cup \{a\} : D \in \mathcal{D}\}$ . Thus the inequality

$$(31) \quad |c(\mathcal{A})| \geq |c(\mathcal{A})| + |c(\mathcal{C})|$$

holds. However,  $\mathcal{B}$  is a system in  $H - \{e\}$ , we may apply the induction hypothesis for  $h - 1$

$$(32) \quad |c(\mathcal{B})| \geq F_l(n - n_2).$$

Further, applying the induction hypothesis for  $l - 1$  we obtain

$$(33) \quad |c(\mathcal{C})| \geq F_{l-1}(n_2).$$

It follows from (31), (32) and (33) that

$$(34) \quad F_l(n - n_2) + F_{l-1}(n_2) \leq |c(\mathcal{A})|.$$

Using the result of Point 5, inequality (5) follows from (34) and (7) by (30), since (7) is proved already for  $l$ .

7. Now prove statement (6) for  $l$  by induction over  $h$ . If  $h = l$ , it is trivial. Assume we have proved (6) for all sets  $|G| = |H| < h$ , and prove for  $h$ . The proof will be similar to the proof of the previous point.

Let  $\mathcal{A}_1$  and  $\mathcal{A}_2$  be given by

$$\mathcal{A}_1 = \{A : A \in \mathcal{A}, A \subset G\},$$

and

$$\mathcal{A}_2 = \{A : A \in \mathcal{A}, A \subset H\}.$$

If  $|\mathcal{A}_1| = r$  and  $|\mathcal{A}_2| = s$ , there are two elements  $e \in G$  and  $f \in H$ , such that  $e$  is contained by at most  $\frac{r \cdot l}{h}$ , and  $f$  is contained by at most  $\frac{s \cdot l}{h}$  sets  $A_i$ . Define the following systems:

$$\mathcal{B} = \{A : A \in \mathcal{A}, e \notin A, f \notin A\},$$

$$\mathcal{C}_1 = \{A - \{e\} : A \in \mathcal{A}_1, e \in A\}$$

and

$$\mathcal{C}_2 = \{A - \{f\} : A \in \mathcal{A}_2, f \in A\},$$

where

$$(35) \quad r_2 = |\mathcal{C}_1| \leq \frac{r \cdot l}{h}$$

and

$$(36) \quad s_2 = |\mathcal{C}_2| \leq \frac{s \cdot l}{h}.$$

Naturally,

$$c(\mathcal{B}) \subset c(\mathcal{A}),$$

$$c(\mathcal{C}_1) \cup e \subset c(\mathcal{A})$$

and

$$c(\mathcal{C}_2) \cup f \subset c(\mathcal{A}).$$

Thus the inequality

$$(37) \quad |c(\mathcal{A})| \geq |c(\mathcal{B})| + |c(\mathcal{C}_1)| + |c(\mathcal{C}_2)| = |c(\mathcal{B})| + |c(\mathcal{C}_1 \cup \mathcal{C}_2)|$$

holds. However  $\mathcal{B}$  is a system in  $G \cup H - \{e\} - \{f\}$ , we may apply our induction hypothesis for  $h - 1$ :

$$(38) \quad |c(\mathcal{B})| \geq \binom{h-1}{l-1} + F_l \left( n - r_2 - s_2 - \binom{h-1}{l} \right).$$

Further, applying the induction hypothesis for  $l - 1$  we obtain

$$(39) \quad |c(\mathcal{C}_1 \cup \mathcal{C}_2)| \geq \binom{h-1}{l-2} + F_{l-1} \left( r_2 + s_2 - \binom{h-1}{l-1} \right).$$

It follows from (37), (38) and (39) that

$$(40) \quad \binom{h}{l-1} + F_l \left( n - r_2 - s_2 - \binom{h-1}{l} \right) + F_{l-1} \left( r_2 + s_2 - \binom{h-1}{l-1} \right) \leq |c(\mathcal{A})|.$$

Now we should like to use inequality (7) which is valid under condition (25) (Point 5). For this reason we have to verify only

$$(41) \quad r_2 + s_2 - \binom{h-1}{l-1} \leq \frac{\left[ n - r_2 - s_2 - \binom{h-1}{l} + r_2 + s_2 - \binom{h-1}{l-1} \right] \cdot l}{a_l^* \left( n - \binom{h}{l} \right)} = \frac{\left[ n - \binom{h}{l} \right] \cdot l}{a_l^* \left( n - \binom{h}{l} \right)}.$$

However

$$(42) \quad r_2 + s_2 - \binom{h-1}{l-1} \leq \frac{\left[ r + s - \binom{h}{l} \right] \cdot l}{h} = \frac{\left[ n - \binom{h}{l} \right] \cdot l}{h}$$

is an immediate consequence of (35) and (36). Since  $n \leq 2 \binom{h}{l}$  is a condition of Theorem 2,  $a_l^* \left( n - \binom{h}{l} \right) \leq h$  holds and (42) results (41). Thus we can use

(7) for this case:

$$(43) \quad F_l \left( n - \binom{h}{l} \right) \leq F_l \left( n - r_2 - s_2 - \binom{h-1}{l} \right) + F_{l-1} \left( r_2 + s_2 - \binom{h-1}{l-1} \right).$$

Finally, (40) and (43) gives the desired inequality, and the whole proof is finished.

Now we consider a natural generalization of the problem of Theorem 1. The problem is to determine the minimum of  $|c^k(\mathcal{A})|$ , where  $1 \leq k \leq l$ ,  $c^k(\mathcal{A}) = c(c^{k-1}(\mathcal{A}))$  and  $c^1(\mathcal{A}) = c(\mathcal{A})$ . It is not difficult to conjecture what is the result. To the theorem we need the following notation:

$$F_l^k(n) = \binom{a_l(n, l)}{l-k} + \binom{a_{l-1}(n, l)}{l-1-k} + \dots + \binom{a_{l(n, l)}(n, l)}{t(n, l) - k} \quad (1 \leq k \leq l),$$

where  $\binom{a}{b} = 0$  if  $b < 0$ .

**THEOREM 3.** Let  $h, n, l$  and  $k$  be given integers with the properties

$$h \geq 1, \quad 1 \leq k \leq l \leq h \quad \text{and} \quad 1 \leq n \leq \binom{h}{l}.$$

If  $H$  is a set of  $h$  elements and

$$\mathcal{A} = \{A_1, \dots, A_n\}, \quad |A_i| = l \quad (i = 1, \dots, n)$$

a system of different subsets of  $H$ , then

$$\min |c^k(\mathcal{A})| = F_l^k(n),$$

where the minimum runs over all such systems  $\mathcal{A}$ .

**PROOF.** It is easy to see by induction over  $l$ , that  $|c^k(\mathcal{M}(h, n, l))| = F_l^k(n)$ . Thus, we have to prove only

$$(44) \quad |c^k(\mathcal{A})| \geq F_l^k(n).$$

This will be proved by induction over  $k$ . For  $k = 1$  Theorem 3 gives Theorem 1. Assume now (44) is true for values smaller than  $k$ , and prove for  $k$ . Obviously,

$$c^k(\mathcal{A}) = c(c^{k-1}(\mathcal{A}))$$

holds and using the induction hypothesis and Theorem 1 we obtain

$$(45) \quad |c^k(\mathcal{A})| \geq F_{l-(k-1)}(F_l^{k-1}(n)).$$

(a) If  $t(n, l) - (k-1) > 0$ , then

$$F_l^{k-1}(n) = \binom{a_l(n, l)}{l-(k-1)} + \dots + \binom{a_{l(n, l)}(n, l)}{t(n, l) - (k-1)}$$

is an expression of type (1). That is

$$(46) \quad \begin{aligned} t(F_l^{k-1}(n), l - k + 1) &= t(n, l) - k + 1 \\ a_i(F_l^{k-1}(n), l - k + 1) &= a_{i+k-1}(n, l) \quad (t(n, l) - k + 1 \leq i \leq l - k + 1) \end{aligned}$$

and

$$(47) \quad \begin{aligned} F_{l-k+1}(F_l^{k-1}(n)) &= \sum_{i=t(n, l)-k+1}^{l-k+1} \binom{a_i(F_l^{k-1}(n), l - k + 1)}{i - 1} = \\ &= \sum_{i=t(n, l)-k+1}^{l-k+1} \binom{a_{i+k-1}(n, l)}{i - 1} = \sum_{j=i(n, l)}^l \binom{a_j(n, l)}{j - k} = F_l^k(n), \end{aligned}$$

which proves (44) and (45).

(b) If  $t(n, l) - k + 1 \leq 0$ , then (46) does not hold. However in this case

$$F_l^{k-1}(n) - 1 = \binom{a_1(n, l)}{l - k + 1} + \dots + \binom{a_k(n, l)}{1}$$

and

$$\begin{aligned} t(F_l^{k-1}(n) - 1, l - k + 1) &= 1 \\ a_i(F_l^{k-1}(n) - 1, l - k + 1) &= a_{i+k-1}(n, l) \quad (1 \leq i \leq l - k + 1) \end{aligned}$$

hold. Further, the equation

$$(48) \quad \begin{aligned} F_{l-k+1}(F_l^{k-1}(n) - 1) &= \sum_{i=1}^{l-k+1} \binom{a_i(F_l^{k-1}(n) - 1, l - k + 1)}{i - 1} = \\ &= \sum_{i=1}^{l-k+1} \binom{a_{i+k-1}(n, l)}{i - 1} = \sum_{j=k}^l \binom{a_j(n, l)}{j - k} = \sum_{j=i(n, l)}^l \binom{a_j(n)}{j - k} = F_l^k(n) \end{aligned}$$

is true in this case instead of (47). If we prove

$$(49) \quad F_{l-k+1}(F_l^{k-1}(n)) = F_{l-k+1}(F_l^{k-1}(n) - 1),$$

then (44) follows from (45), (49) and (48). (49) will be proved by the following simple lemma.

LEMMA 2. If  $t(m, r) = 1$ , then

$$F_r(m + 1) = F_r(m).$$

PROOF. Let  $s$  be the least index such that  $a_s(m, r) > a_{s-1}(m, r) + 1$  ( $2 \leq s \leq r$ ). If there is not such  $s$ , let  $s$  be equal to  $r + 1$ . Thus, we can write

$$\begin{aligned} m &= \binom{a_r(m, r)}{r} + \dots + \binom{a_{s-1}(m, r)}{s - 1} + \binom{a_{s-1}(m, r) - 1}{s - 2} + \dots + \\ &\quad + \dots + \binom{a_{s-1}(m, r) - (s - 2)}{1} \end{aligned}$$

and

$$m + 1 = \binom{a_r(m, r)}{r} + \dots + \binom{a_{s-1}(m, r) + 1}{s-1}.$$

Now it is not difficult to see, that

$$\begin{aligned} F_r(m) &= \binom{a_r(m, r)}{r-1} + \dots + \binom{a_{s-1}(m, r)}{s-2} + \binom{a_{s-1}(m, r) - 1}{s-3} + \dots + \\ &+ \dots + \binom{a_{s-1}(m, r) - (s-2)}{0} = \binom{a_r(m, r)}{r-1} + \dots + \binom{a_{s-1}(m, r) + 1}{s-2} = \\ &= F_r(m + 1), \end{aligned}$$

which proves the lemma and Theorem 3.

### § 3. Solution of an Erdős-problem

Let  $H$  be a finite set of  $h$  elements, and  $\mathcal{A}$  a system of subsets of  $H$ :

$$\mathcal{A} = \{A_1, A_2, \dots, A_n\}, A_i \subset H, |A_i| = l \quad (1 \leq i \leq n).$$

ERDŐS proposed the following problem. For which numbers  $n$  can we construct a system  $\mathcal{B}$  with the properties

$$\mathcal{B} = \{B_1, B_2, \dots, B_n\}, B_i \subset A_i, |B_i| = l - k \quad (1 \leq i \leq n).$$

In the solution we use the well-known marriage problem. It is clear in this connection, that it is a very important question, in which cases does  $F_l^k(n) < n$ ,  $F_l^k(n) = n$  or  $F_l^k(n) > n$  hold. The following sequence of lemmas deals with this problem.

LEMMA 3. *If  $1 \leq k \leq l$  and  $x$  are positive integers, then*

$$f(x) = \binom{x}{l-k} - \binom{x}{l}$$

*is a monotone increasing function between  $l$  and  $2l - k - 2$  but it is a monotone decreasing function from  $2l - k - 1$ . The values  $f(2l - k - 2)$  and  $f(2l - k - 1)$  are equal.*

PROOF. Let  $0 \leq a < b \leq x$  be integers. It is easy to see that  $\binom{x}{a} - \binom{x}{b} < 0$ ,  $\binom{x}{a} - \binom{x}{b} = 0$  and  $\binom{x}{a} - \binom{x}{b} > 0$ , respectively, if  $a + b < x$ ,  $a + b = x$  and  $a + b > x$ , respectively.

Consider the difference  $f(x+1) - f(x) = \binom{x}{l-k-1} - \binom{x}{l-1}$ . Using the above remark we obtain that

$$f(x+1) - f(x) < 0 \quad \text{if} \quad 2l - k - 2 < x,$$

$$f(x+1) - f(x) = 0 \quad \text{if} \quad 2l - k - 2 = x,$$

and finally,

$$f(x+1) - f(x) > 0 \quad \text{if} \quad 2l - k - 2 > x.$$

This completes the proof.

The following two lemmas are immediate consequences of Lemma 3.

LEMMA 3a. *If  $1 \leq k \leq l$  and  $x$  are positive integers, then*

$$\binom{x}{l-k} - \binom{x}{l} \leq \binom{2l-k-1}{l-k} - \binom{2l-k-1}{l}.$$

LEMMA 3b. *If  $1 \leq k \leq l$  and  $x \geq 2l - k + 1$  are positive integers, then*

$$\binom{x}{l-k} - \binom{x}{l} \leq \binom{2l-k+1}{l-k} - \binom{2l-k+1}{l}.$$

LEMMA 4. *If  $1 \leq k \leq m$ , then*

$$\sum_{i=k}^{m-1} \left[ \binom{2i-k-1}{i-k} - \binom{2i-k-1}{i} \right] \leq \binom{2m-k-1}{m-k} - \binom{2m-k-1}{m}.$$

PROOF. Let  $a$  and  $b$  be positive integers, where  $\frac{a}{2} \leq b \leq a-1$ . Then

$$(50) \quad \binom{a}{b} - \binom{a}{b+1} = \binom{a}{b} \left[ 1 - \frac{a-b}{b+1} \right] = \binom{a}{b} \left[ \frac{2b-a+1}{b+1} \right],$$

and similarly

$$(51) \quad \binom{a+2}{b+1} - \binom{a+2}{b+2} = \binom{a+2}{b+1} \left[ 1 - \frac{a-b+1}{b+2} \right] = \binom{a+2}{b+1} \left[ \frac{2b-a+1}{b+2} \right].$$

Further

$$(52) \quad \binom{a+2}{b+1} \left[ \frac{2b-a+1}{b+2} \right] = \binom{a}{b} \left[ \frac{2b-a+1}{b+1} \right] \cdot \left[ \frac{(a+2)(a+1)}{(b+2)(a-b+1)} \right],$$

where

$$\frac{a+1}{b+2} \geq 1,$$

and

$$\frac{a+2}{a-b+1} \geq \frac{a+2}{\frac{a}{2}+1} = 2.$$

That is

$$(53) \quad \binom{a+2}{b+1} - \binom{a+2}{b+2} \geq 2 \left[ \binom{a}{b} - \binom{a}{b+1} \right]$$

follows from (50), (51) and (52).

Applying (53) for  $a = 2i = k - 1$ , and  $b = i - 1$ , ( $1 \leq k \leq i$ ), we obtain

$$\binom{2(i+1)-k-1}{i} - \binom{2(i+1)-k-1}{i+1} \geq 2 \left[ \binom{2i-k-1}{i-1} - \binom{2i-k-1}{i} \right],$$

or

$$(54) \quad \binom{2(i+1)-k-1}{i+1-k} - \binom{2(i+1)-k-1}{i+1} \geq 2 \left[ \binom{2i-k-1}{i-k} - \binom{2i-k-1}{i} \right].$$

Prove now the lemma by induction over  $m$ . If  $m = k$ , the statement is trivial. Let the lemma be true for  $m$  and prove it for  $m + 1$ .

$$\begin{aligned} \sum_{i=k}^m \left[ \binom{2i-k-1}{i-k} - \binom{2i-k-1}{i} \right] &= \sum_{i=k}^{m-1} \left[ \binom{2i-k-1}{i-k} - \binom{2i-k-1}{i} \right] + \\ &\quad + \binom{2m-k-1}{m-k} - \binom{2m-k-1}{m} \end{aligned}$$

and by induction hypothesis and (54)

$$\begin{aligned} \sum_{i=k}^m \left[ \binom{2i-k-1}{i-k} - \binom{2i-k-1}{i} \right] &\leq 2 \left[ \binom{2m-k-1}{m-k} - \binom{2m-k-1}{m} \right] \leq \\ &\leq \binom{2(m+1)-k-1}{m+1-k} - \binom{2(m+1)-k-1}{m+1} \end{aligned}$$

holds, which proves Lemma 4.

LEMMA 5. If  $1 \leq k \leq l$  and  $2l - k < a_l(n, l)$  then

$$(55) \quad F_l^k(n) < n.$$

PROOF. We may use Lemma 3b:

$$(56) \quad \binom{a_l(n, l)}{l-k} - \binom{a_l(n, l)}{l} \leq \binom{2l-k+1}{l-k} - \binom{2l-k+1}{l}.$$

On the other hand, by Lemma 3a

$$\binom{a_l(n, l)}{i-k} - \binom{a_l(n, l)}{i} \leq \binom{2i-k-1}{i-k} - \binom{2i-k-1}{i} \quad (k \leq i \leq l-1)$$

holds and summarizing it we obtain

$$\sum_{i=k}^{l-1} \left[ \binom{a_i(n, l)}{i-k} - \binom{a_i(n, l)}{i} \right] \leq \sum_{i=k}^{l-1} \left[ \binom{2i-k-1}{i-k} - \binom{2i-k-1}{i} \right].$$

Applying now Lemma 4 and (54):

$$\begin{aligned} \sum_{i=k}^{l-1} \left[ \binom{a_i(n, l)}{i-k} - \binom{a_i(n, l)}{i} \right] &\leq \binom{2l-k-1}{l-k} - \binom{2l-k-1}{l} < \\ &< \binom{2(l+1)-k-1}{l+1-k} - \binom{2(l+1)-k-1}{l+1}. \end{aligned}$$

Obviously,

$$(57) \quad \sum_{i=1}^{l-1} \left[ \binom{a_i(n, l)}{i-k} - \binom{a_i(n, l)}{i} \right] < \binom{2(l+1)-k-1}{l+1-k} - \binom{2(l+1)-k-1}{l+1}$$

also holds, since we added a nonpositive number to the left side. If we sum (56) and (57) the obtained inequality

$$\begin{aligned} F_l^k(n) - n &< \binom{2l-k+1}{l-k} - \binom{2l-k+1}{l} + \binom{2(l+1)-k-1}{l+1-k} - \\ &\quad - \binom{2(l+1)-k-1}{l+1} = 0 \end{aligned}$$

results (55).

LEMMA 6. *If  $1 \leq k \leq l$  and  $2l - k > a_l(n, l)$  then*

$$(58) \quad F_l^k(n) > n.$$

PROOF. We know that

$$(59) \quad a_i(n, l) \leq a_l(n, l) - (l - i).$$

If  $l > i \geq a_l(n, l) - (l - k)$ , then  $a_i(n, l) - (l - i) \leq 2i - k$  and by (59)

$$a_i(n, l) \leq 2i - k$$

holds. In this case, obviously

$$(60) \quad \binom{a_i(n, l)}{i-k} - \binom{a_i(n, l)}{i} \geq 0$$

follows. If  $k \leq i < a_l(n, l) - (l - k)$ , then by (59) and Lemma 3

$$(61) \quad \binom{a_i(n, l)}{i-k} - \binom{a_i(n, l)}{i} \geq \binom{a_i(n, l) - (l - i)}{i-k} - \binom{a_i(n, l) - (l - i)}{i}$$

holds, but it is trivially true for  $i < k$ , too.

Sum (60) and (61)

$$\sum_{i=1}^{l-1} \left[ \binom{a_i(n, l)}{i-k} - \binom{a_i(n, l)}{i} \right] \geq \sum_{i=1}^{a_i(n, l)-l+k-1} \left[ \binom{a_i(n, l) - (l-i)}{i-k} - \binom{a_i(n, l) - (l-i)}{i} \right] = \binom{2a_i(n, l) - 2l + k}{a_i(n, l) - l - 1} - \binom{2a_i(n, l) - 2l + k}{a_i(n, l) - l + k - 1} + 1.$$

That is

$$(62) \quad F_l^k(n) - n \geq \binom{2a_i(n, l) - 2l + k}{a_i(n, l) - l - 1} - \binom{2a_i(n, l) - 2l + k}{a_i(n, l) - l + k - 1} + 1 + \binom{a_i(n, l)}{l-k} - \binom{a_i(n, l)}{l}$$

s true. Here

$$(63) \quad \binom{2a_i(n, l) - 2l + k}{a_i(n, l) - l - 1} - \binom{2a_i(n, l) - 2l + k}{a_i(n, l) - l + k - 1} \geq \binom{a_i(n, l) - 1}{a_i(n, l) - l - 1} - \binom{a_i(n, l) - 1}{a_i(n, l) - l + k - 1}$$

because of Lemma 3. However we can write the right hand side of (63) in the form

$$(64) \quad \binom{a_i(n, l) - 1}{a_i(n, l) - l - 1} - \binom{a_i(n, l) - 1}{a_i(n, l) - l + k - 1} = \binom{a_i(n, l) - 1}{l} - \binom{a_i(n, l) - 1}{l-k}.$$

Here

$$\binom{a_i(n, l) - 1}{l} - \binom{a_i(n, l) - 1}{l-k} = \left[ \binom{a_i(n, l)}{l} - \binom{a_i(n, l)}{l-k} \right] - \left[ \binom{a_i(n, l) - 1}{l-1} - \binom{a_i(n, l) - 1}{l-k-1} \right]$$

and since  $2l - k - 1 > a_i(n, l) - 1$  by supposition of the lemma, thus

$$(65) \quad \binom{a_i(n, l) - 1}{l} - \binom{a_i(n, l) - 1}{l-k} \geq \binom{a_i(n, l)}{l} - \binom{a_i(n, l)}{l-k}.$$

Finally, (65), (64) (63) and (62) give

$$F_l^k(n) - n \geq 1$$

which proves our lemma.

LEMMA 7. *If*

$$(66) \quad n > \binom{2l-k}{l} + \binom{2(l-1)-k}{l-1} + \dots + \binom{k}{k},$$

then

$$F_l^k(n) < n.$$

On the other hand, if

$$(67) \quad n \leq \binom{2l-k}{l} + \binom{2(l-1)-k}{l-1} + \dots + \binom{k}{k},$$

then

$$F_l^k(n) \geq n$$

with equality only if

$$(68) \quad n = \binom{2l-k}{l} + \binom{2(l-1)-k}{l-1} + \dots + \binom{2s-k}{s}$$

for some  $s$  ( $k \leq s \leq l$ ).

PROOF. Consider first the case of (66). If  $a_i(n, l) = 2i - k$  ( $k \leq i \leq l$ ), then  $t(n, l) < k$  and

$$n = \binom{2l-k}{l} + \dots + \binom{k}{k} + \binom{k-1}{k-1} + \dots + \binom{t(n, l)}{t(n, l)}.$$

Obviously,

$$F_l^k(n) = \binom{2l-k}{l} + \dots + \binom{k}{k},$$

thus  $F_l^k(n) < n$  holds.

In the contrary case

$$a_r(n, l) > 2r - k$$

$$a_i(n, l) = 2i - k \quad (r < i \leq l)$$

hold for some  $r$  ( $k \leq r < l$ ). Since

$$n - F_l^k(n) = \left[ \binom{a_r(n, l)}{r} + \dots \right] - \left[ \binom{a_r(n, l)}{r-k} + \dots \right],$$

the statement follows by Lemma 5:

$$\binom{a_r(n, l)}{r} + \dots > F_r^k \left[ \binom{a_r(n, l)}{r} + \dots \right].$$

The case (67) may occur in two different ways.

1. If (68) holds, then obviously  $F_l^k(n) = n$

2. For some  $r$  ( $k < r \leq l$ ),

$$a_r(n, l) < 2r - k,$$

and

$$a_i(n, l) - 2i - k \quad (r < i \leq l).$$

Since

$$n - F_l^k(n) = \left[ \binom{a_r(n, l)}{r} + \dots \right] - \left[ \binom{a_r(n, l)}{r-k} + \dots \right],$$

the statement follows by Lemma 6:

$$\binom{a_r(n, l)}{r} + \dots < F_r^k \left[ \binom{a_r(n, l)}{r} + \dots \right].$$

**THEOREM 4.** *Let  $1 \leq k \leq l \leq h$  be positive integers,  $H$  a set of  $h$  element and*

$$\mathcal{A} = \{A_1, \dots, A_n\}, \quad |A_i| = l \quad (1 \leq i \leq n)$$

*a system of subsets of  $H$ . If*

$$(69) \quad n \leq \binom{2l-k}{l} + \binom{2(l-1)-k}{l-1} + \dots + \binom{k}{k},$$

*there exists a system*

$$(70) \quad \mathcal{B} = \{B_1, \dots, B_n\}, \quad |B_i| = l - k, \quad B_i \subset A_i \quad (1 \leq i \leq n)$$

*but in the case of*

$$(71) \quad n > \binom{2l-k}{l} + \binom{2(l-1)-k}{l-1} + \dots + \binom{k}{k}$$

*not necessarily.*

**PROOF.** First we prove the latter case. If (71) holds then by Lemma 7  $F_l^k(n) < n$ . We know (Theorem 1) that there exists a system  $\mathcal{A}$  such that  $|c^k(\mathcal{A})| = F_l^k(n)$ . Thus, a system  $\mathcal{B}$  satisfying (70) does not exist.

In the proof of the existence of  $\mathcal{B}$  in the case of (69) we use the well-known marriage problem [2]:

**THEOREM OF ORE.** *Let  $E$  and  $F$  be disjoint sets and  $G$  a graph on  $E \cup F$ . Assume  $G$  has the property that for arbitrary  $D \subset E$  there is a set  $H \subset F$  such that every element of  $H$  is connected with at least one element of  $D$  and  $|H| \geq |D|$ . Then there exists a one-to-one mapping between  $E$  and a subset  $K$  of  $F$ , such that the associating vertices are connected in  $G$ .*

In our case  $E = \mathcal{A}$ ,  $F = c^k(\mathcal{A})$  and  $A \in \mathcal{A}$ ,  $B \in c^k(\mathcal{A})$  are connected if and only if  $A \supset B$ . Thus, it is sufficient to verify that for every subsystem

$$\mathcal{C} = \{A_{i_1}, \dots, A_{i_m}\} \subset \mathcal{A}$$

there are at least  $m$  sets in  $c^k(\mathcal{A})$ , which are contained in one of  $A_{i_j}$  ( $1 \leq j \leq m$ ). However,  $m \leq n$ , thus by (69)

$$m \leq \binom{2l-k}{l} + \dots + \binom{k}{k}$$

and Lemma 7 gives

$$(72) \quad F_l^k(m) \geq m.$$

Use now Theorem 1:

$$|c^k(\mathcal{C})| \geq F_l^k(m).$$

This and (72) results  $|c^k(\mathcal{C})| \geq m$ , which means that our graph has the property prescribed in the used theorem. Applying the theorem the obtained one-to-one mapping gives just the desired system  $\mathcal{B}$ .

COROLLARY. If  $2l - k \geq h$ , then (69) always holds and a system  $\mathcal{B}$  satisfying (70) always exists.

This is an immediate consequence of the inequality

$$\binom{h}{l} \leq \binom{2l-k}{l} \leq \binom{2l-k}{l} + \dots + \binom{k}{k}$$

and the fact that  $\mathcal{A}$  has at most  $\binom{h}{l}$  elements.

#### REFERENCES

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