

ON TRIGONOMETRIC SUMS WITH GAPS

by
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A well known theorem states as follows:¹

Let $n_1 < n_2 < \dots, n_{k+1}/n_k > A > 1$ be an infinite sequence of real numbers and $\sum_{k=1}^{\infty} (a_k^2 + b_k^2)$ a divergent series satisfying

$$(1) \quad \lim_{N=\infty} (a_N^2 + b_N^2)^{1/2} \left(\sum_{k=1}^N a_k^2 + b_k^2 \right)^{-1/2} = 0.$$

Then

$$(2) \quad \lim_{n \rightarrow \infty} \left| \int_t \left\{ \sum_{k=1}^N (a_k \cos 2\pi n_k t + b_k \sin 2\pi n_k t) < \omega \left(\frac{1}{2} \sum_{k=1}^N (a_k^2 + b_k^2) \right)^{1/2} \right\} \right| = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\omega} e^{-u^2/2} du.$$

($\int_t \{ \quad \}$ denotes the Lebesgue measure of the set in question).

In the present paper I shall weaken the lacunarity condition $n_{k+1}/n_k > A > 1$. In fact I shall prove the following

Theorem 1. *Let $n_1 < n_2 < \dots$ be an infinite sequence of integers satisfying*

$$(3) \quad n_{k+1} > n_k \left(1 + \frac{c_k}{k^{1/2}} \right)$$

where $c_k \rightarrow \infty$. Then

$$(4) \quad \lim_{N=\infty} \left| \int_t \left\{ \sum_{k=1}^N \left(\cos 2\pi n_k (t - \vartheta_k) < \omega \cdot N^{1/2} \right) \right\} \right| = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\omega} e^{-u^2/2} du.$$

It seems likely that the Theorem remains true if it is not assumed that the n_k are integers. On the other hand if $n_{k+1}/n_k \rightarrow 1$ is an arbitrary sequence of integers it is easy to construct examples which show that (1) is not enough

¹ R. SALEM and A. ZYGMUND: "On lacunary trigonometric series I. and II.," *Proc. Math. Acad. Sci. USA* **33** (1947) 333-338 and **34** (1948) 54-62.

For the history of the problem see M. KAC: "Probability methods in analysis and number theory", *Bull. Amer. Math. Soc.* **55** (1949) 641-665.

for the truth of (2). It is possible that (3) and

$$\lim_{N=\infty} \left(\sum' a_k^2 + b_k^2 \right)^{\frac{1}{2}} \left(\sum_{k=1}^N (a_k^2 + b_k^2) \right)^{-\frac{1}{2}} = 0$$

where in $\sum' \frac{1}{2} n_N < n_k \leq n_N$ suffices for the truth of our Theorem. But I can not at present decide this question and in this paper only consider the case $a_k = b_k = 1$.

I can show that Theorem 1. is best possible in the following sense: To every constant c there exists a sequence n_k for which $n_{k+1} > n_k \left(1 + \frac{c}{k^{\frac{1}{2}}} \right)$ but (4) is not true. To see this let u_k tend to infinity sufficiently fast. Put

$$n_{k^2+l} = n_k + lc_1 \left[\frac{n_k}{k} \right], \quad 1 \leq l \leq 2k+1.$$

Clearly $n_{r+1} > n_r \left(1 + \frac{c}{r^{\frac{1}{2}}} \right)$ if c_1 is sufficiently large and it is not difficult to see that (4) can not be satisfied. We do not give the details.

Further I can prove the following

Theorem 2. *Let $n_1 < n_2 < \dots$ be an infinite sequence of integers for which for every $\varepsilon > 0$ there exists a $k_0 = k_0(\varepsilon)$ so that for every $k > k_0$*

$$(5) \quad n_{k+1} > n_k + n_{k - [\varepsilon k^{1/2}]}.$$

Then (4) holds.

It is not difficult to construct sequences for which (3) holds but (5) does not hold and sequences for which (5) holds and (3) not, or Theorems 1 and 2 are incomparable. (3) seems to be easier to verify, thus Theorem 1 is probably more useful. We will not give the proof of Theorem 2 since it is similar to that of Theorem 1.

To simplify the computations we will work out the proof of Theorem 1 only for a cosine series, the proof of the general case follows the same lines.

Theorem 1'. *Let $n_1 < n_2 < \dots$ be an infinite sequence of integers satisfying (3). Then*

$$(4') \quad \lim_{N=\infty} \left| E \left\{ \sum_{k=1}^N \cos 2\pi n_k t < \omega \left(\frac{N}{2} \right)^{\frac{1}{2}} \right\} \right| = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\omega} e^{-u^2/2} du.$$

A well known theorem of Chebyshev implies that to prove Theorem 1' it will suffice to show that for every l , $1 \leq l < \infty$

$$(6) \quad \lim_{N=\infty} I_N^{(l)} = \lim_{N=\infty} \int_0^1 \left(\frac{\sum_{k=1}^N \cos 2\pi n_k t}{\left(\frac{N}{2} \right)^{\frac{1}{2}}} \right)^l dt = (2\pi)^{-\frac{1}{2}} \int_{-\infty}^{\infty} x^l e^{-x^2/2} dx = \begin{cases} 0 & \text{if } l \text{ is odd,} \\ \frac{l!}{2^{l/2} \left(\frac{l}{2} \right)!} & \text{if } l \text{ is even.} \end{cases}$$

It is easy to see that $(\varepsilon_i = \pm 1, 1 \leq i \leq l)$

$$(7) \quad \int_0^1 \prod_{i=1}^l \cos 2\pi n_i t = \frac{1}{2^l} \int_0^1 \sum_{\varepsilon_1, \dots, \varepsilon_l} \cos \left(2\pi \sum_{i=1}^l \varepsilon_i n_i t \right) dt = \frac{h(n_1, \dots, n_l)}{2^l}$$

where $h(n_1, \dots, n_l)$ denotes the number of solutions of $\sum_{i=1}^l \varepsilon_i n_i = 0$. From (7) we have

$$(8) \quad \left(\frac{N}{2} \right)^{l/2} I_N^{(l)} = \frac{1}{2^l} \sum h(n_{i_1}, \dots, n_{i_l})$$

where i_1, \dots, i_l runs through all the l -tuples formed from the integers $1 \leq r \leq N$ (where order counts). Clearly $\sum h(n_{i_1}, \dots, n_{i_l})$ equals the number of solutions of

$$(9) \quad \sum_{i=1}^l \varepsilon_i n_{r_i} = 0, \quad 1 \leq r_i \leq N \quad (\text{order counts here too}).$$

Thus to estimate $I_N^{(l)}$ we only have to estimate the number of solutions of (9). Assume first l even $l = 2s$. Then (9) has trivial solutions such that among the terms in (9) each n_r occurs the same number of times with a positive as with a negative sign. The number of these trivial solutions clearly equals

$$(10) \quad (1 + o(1)) \frac{l!}{\left(\frac{l}{2}\right)!} N^{l/2}.$$

Now we prove the following

Lemma 1. Let $\{n_k\}$ be a sequence of integers satisfying (3). Denote by $g_l(A, N)$ the number of solutions of

$$(11) \quad \sum_{i=1}^l \varepsilon_i n_{r_i} = A, \quad 1 \leq r_1 \leq \dots \leq r_l \leq N$$

where the trivial solutions are excluded.

Then

$$(12) \quad \max_A g_l(A, N) = o(N^{l/2}).$$

(The trivial solutions can only occur if $A = 0$ and l is even).

From Lemma 1, (8) and (10) it follows that

$$\lim_{N \rightarrow \infty} I_N^{(l)} = \begin{cases} 0 & \text{if } l \text{ is odd} \\ \frac{l!}{2^{l/2} \left(\frac{l}{2}\right)!} & \text{if } l \text{ is even} \end{cases}$$

which implies Theorem 1. Thus to complete our proof it will suffice to prove Lemma 1, and in fact Lemma 1 is the only new and difficult part of our paper.

First we show that the Lemma holds for $l = 1$ and $l = 2$. For $l = 1$ the Lemma is trivial, the number of solutions of (11) is at most one for $l = 1$. Now we need

Lemma 2. *The number of n_i satisfying ($k \rightarrow \infty$)*

$$n_k x^{-1} < n_i < n_k x$$

is $o(k^{1/2} \log x) + o(1) + o((\log x)^2)$.

The Lemma follows immediately from (3). (The term $o(1)$ is needed only for small x and the term $o((\log x)^2)$ only for very large x .)

If $\pm n_{r_1} \pm n_{r_2} = A(n_{r_1} > n_{r_2})$ we must have (by (3))

$$(13) \quad \left| \frac{A}{2} \right| \leq n_{r_1} \leq |A| N.$$

From (13) and Lemma 2 we obtain that the number of solutions of (11) for $l = 2$ is $o(N^{1/2} \log N) = o(N)$ uniformly in A which proves Lemma 1 for $l = 2$.

Now we use induction with respect to l . Assume that (12) holds for all $l' < l$, we shall then prove that (12) holds for l too. We assume now $l \geq 3$ and distinguish four cases.

In case I.

$$(14) \quad \frac{1}{N} n_{r_i} \leq n_{r_{i+1}} \leq n_{r_i}$$

holds for all $1 \leq i \leq l - 1$.

Put $(1 \leq s \leq l - 1)$

$$(15) \quad 2^n \leq \max n_{r_i} / n_{r_{i+1}} = n_{r_s} / n_{r_{s+1}} < 2^{n+1}.$$

Clearly $0 \leq n \leq \log N / \log 2$. Evidently there are at most N choices for n_{r_1} . Let $i < l - 1$. If n_{r_1}, \dots, n_{r_i} have already been determined then by (15) and Lemma 2 there are at most $o(N^{1/2} n)$ choices for $n_{r_{i+1}}$. Now we show that for n_{r_s} there at most are $o(N^{1/2} / 2^n) + o(1) = o\left(\frac{N^{1/2}}{2^{n/4}}\right)$ choices (if $n_{r_1}, \dots, n_{r_{s-1}}$ has already been chosen). To see this observe that from (15) we have

$$(16) \quad \left| \sum_{i=s+1}^l \varepsilon_i n_{r_i} \right| < \frac{l \cdot n_{r_s}}{2^n}.$$

Thus from (11) and (16)

$$(17) \quad A - \sum_{i=1}^{s-1} \varepsilon_i n_i = \varepsilon_s n_{r_s} + \frac{\theta l}{2^n} n_{r_s}, \quad |\theta| < 1.$$

(17) implies that n_{r_s} must lie in an interval (α, β) with $\alpha < \beta < \alpha \left(1 + \frac{cl}{2^n}\right)$.

Thus from Lemma 2 there are at most $o\left(\frac{N^{1/2}}{2^n}\right) + o(1) = o\left(\frac{N^{1/2}}{2^{n/4}}\right)$ choices for n_{r_s} as stated. Finally if $n_{r_1}, \dots, n_{r_{l-1}}$ has already been determined there are at most 2^{l-1} choices for n_{r_l} (i. e. $\sum_{i=1}^{l-1} \varepsilon_i n_{r_i}$ can be chosen in 2^{l-1} ways). Thus the

total number of choices for n_{r_1}, \dots, n_{r_l} satisfying (15) is at most

$$(18) \quad cN(o(N^{1/2}n))^{l-3} o\left(\frac{N^{1/2}}{2^{n/4}}\right) = o(N^{l/2}) \frac{n^{l-3}}{2^{n/4}}.$$

From (18) we evidently obtain that the number of solutions of (11) in case I is

$$o(N^{l/2}) \sum_{n=0}^{\infty} \frac{n^{l-3}}{2^{n/4}} = o(N^{l/2}).$$

In case II (14) holds for $i < j, j \geq 3$ and for $i = j \leq l - 1$

$$(19) \quad n_{r_{j+1}} < \frac{1}{N} n_{r_j}.$$

We show that if $n_{r_1}, \dots, n_{r_{j-1}}$ has already been determined, then there are only a bounded number of choices of n_{r_j} . To see this observe that by (19)

$$\left(\sum_{i=j+1}^l \varepsilon_i n_{r_i} \right) < \frac{l}{N} n_{r_j}.$$

Thus from (11)

$$(20) \quad A - \sum_{i=1}^{j-1} \varepsilon_i n_{r_i} = \varepsilon_j n_{r_j} + \theta \frac{ln_j}{N}, \quad |\theta| < 1$$

or n_{r_j} must lie in an interval (α, β) with $\alpha < \beta < \alpha \left(1 + \frac{cl}{N}\right)$. Thus by Lemma 2 there are only a bounded number of choices for n_{r_j} .

Put

$$(15') \quad 2^n \leq \max_{1 \leq i \leq j-1} n_{r_i}/n_{r_{i+1}} = n_{r_s}/n_{r_{s+1}} < 2^{n+1}.$$

As in case I, there are at most $o(N^{1/2}/2^{n/4})$ choices for n_{r_s} , $o(N^{1/2}n)$ choices for n_{r_i} , $1 < i < j$, $i \neq s$ and at most N choices for n_{r_1} . Thus we see as in case I that for n_{r_1}, \dots, n_{r_j} there are at most $o(N^{j/2})$ choices. If n_{r_1}, \dots, n_{r_j} are already chosen there are 2^j choices for $\sum_{i=1}^j \varepsilon_i n_{r_i}$. Hence there are only $2^j o(N^{j/2}) = o(N^{j/2})$ choices for $\sum_{i=1}^j \varepsilon_i n_{r_i}$. By our induction hypothesis there are $o(N^{(j-1)/2})$ solutions of

$$(21) \quad A - \sum_{i=1}^j \varepsilon_i n_{r_i} = \sum_{i=j+1}^l \varepsilon_i n_{r_i}$$

in $n_{r_{j+1}}, \dots, n_{r_l}$. Thus finally there are $o(N^{l/2})$ solutions of (11) in case II.

In case III (14) holds for $i = 1$, but

$$n_{r_3} < \frac{1}{N} n_{r_2}.$$

The same proof as in case II shows that if n_{r_1} has already been chosen there are only a bounded number of choices for n_{r_2} . Thus since there are at most N choices for n_{r_1} there are at most cn choices for $\varepsilon_1 n_{r_1} + \varepsilon_2 n_{r_2}$. Hence arguing

as in (21) we see that by our induction hypothesis the number of solutions of (11) is $o(N^{1/2})$ in case III too.

In case IV $n_{r_2} < \frac{1}{N} n_{r_1}$ i. e. (14) never holds. We see by the same argument as in (20) that there are only a bounded number of choices for n_{r_1} and therefore again arguing as in (21) we obtain by our induction hypothesis that in case IV (11) has $o(N^{1/2})$ solutions.

Thus combining the four cases we obtain that the number of solutions of (11) is $o(N^{1/2})$ uniformly in A , or (12) — and therefore Lemma 1 is proved. Hence the proof of Theorem 1 is complete.

Let $f(k) \rightarrow \infty$ monotonically. It is easy to see that

$$(22) \quad n_k = [e^{k^{1/2} f(k)}]$$

satisfies (3), hence Theorem 1 holds for the sequence (22).

It is not difficult to see that Lemma 1 is best possible in some sense, namely if (3) is replaced by

$$n_{k+1} > n_k \left(1 + \frac{c}{k^{1/2}}\right) \quad c \text{ independent of } k$$

then (12) in general will not hold. On the other hand (12) may very well hold for special sequences which do not satisfy (3). In particular I would guess that (12) and therefore Theorem 1 will hold if $n_k = [e^{k^\alpha}]$ for every $\alpha > 0$. I cannot even prove this for $\alpha = 1/2$.

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О ЛАКУНАРНЫХ ТРИГОНОМЕТРИЧЕСКИХ РЯДАХ

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Резюме

В работе доказывается следующая теорема: пусть $n < n_2 < \dots$ последовательность натуральных чисел для которых

$$n_{k+1} > n_k \left(1 + \frac{c_k}{\sqrt{k}}\right) \quad (k = 1, 2, \dots),$$

где

$$\lim_{k \rightarrow \infty} c_k = +\infty.$$

Пусть $S_N(t) = \sum_{k=1}^N \cos 2\pi n_k(t - \vartheta_k)$ где вещественные числа ϑ_k произвольные.

Пусть $E_t \{ \}$ обозначает множество тех чисел t в интервале $0 \leq t \leq 1$ для которых условие в скобках выполняется, и пусть $|E_t|$ — мера Lebesgue-а множества E_t . Тогда имеем для всех ω ($-\infty < \omega < \infty$)

$$\lim_{N \rightarrow \infty} |E_t \{S_N(t) < \omega \sqrt{N}\}| = \frac{1}{\sqrt{2\pi}} \int_{-\infty}^{\omega} e^{-u^2/2} du.$$